Temporal HTML

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ABSTRACT

Temporal HTML (THTML) as presented in this thesis is an extension of HTML -- a high level authoring language for World Wide Web documents. THTML incorporates temporal logic into HTML to provide an efficient solution for authoring and maintaining *time-sensitive* web sites. In THTML, the same URL request may result in different HTML pages for different request times. This request time may **be** the local time that the reader sends the request, or it may **be** sometime in the past or even the future. The HTML page sent in response to **the** request is an instance of the page determined by the particular time context, i.e. the reader's request time. The instance is generated using the THTML source files for each web component whose time interval stamps most closely approximate the time specified. With *server push* technology incorporated into the design, the THTML server periodicaily re-instantiates the previously requested **URLs** in each web session and re-issues the updated instances as time elapses. THTML 1.0 implemented the above design partially.

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Chapter 1

Introduction

1.1 Introduction to Temporal HTML

Since Web technology provides an integrated presentation of al1 types of information (text, picture, video and audio) on a single screen, individuals and institutions are encouraged to establish their own Internet presence. This multi-media environment can consist of text, graphies, audio, video, three-dimensional models or any other information that a computer **can** process.

Web technology is based on 2 protocols, namely HTML and HTTP. Based on the TCP/IP Internet layer, HTTP dominates the protocols defined by W3C for the World Wide Web. In HTTP, a set of negotiations is defined between the client and server so that actions performed at the client *cm* **be** interpreted by the server and the correct response can be sent to the client. For example, the *ger* method is defined to retrieve certain resources at the server site. HTML (hypertext markup language) is **normaily** considered a high-level authoring tool to wnte web container documents that hold different types of multimedia information. Web browsers can interpret the different types of data and display them on the screen. HTML is actually a protocol implemented on top of **H?TP** for presenting information on the web. It features **an** *archor* system by which any web page can refer to other pages, without any knowledge of the contents of the other page. It gives the web a better look and feel **than** that of a static electronic library.

Although HTML serves many applications very well as a high-level authoring language for the World Wide Web, it provides no temporal support for a site. With the rapid development of web technology nowadays **and** the appearance of more time sensitive applications on the **web,** it is no longer sufficient to provide only static information. Creating a dynamic site becomes criticai for success. The non-temporal feature of conventiond **I-ïïML** is a serious hindrance to creating a **more** dynarnic World

 \mathbf{I}

Wide Web. Some of the problems are concluded as follows:

- 1. In addition to retrieving the current version, web clients sometimes wish to retrieve historical information. This **occurs** on sites related to the news or stock market where retrievai of information based on time is necessary. However, the existing HTML protocol only supports keeping one version of a document in the system.
- 2. In time critical web applications, such as the stock market, information being viewed at one moment might become stale in the next. HTML does not provide a mechanism for synchronizing the server and clients, so clients will not be informed of changes to the information.
- **3.** The demand for web authors to create **dynamic** rather **than** *staric* sites is increasing dramaticaily. They need to decide **when** items should be posted, and not just *what* should be posted. The existing web servers do not have a mechanism selecting the file corresponding to a given time. The author has to manually update the site to make it **appear** dynamic from the client's perspective.

With the above demands for temporal support on the web, a lot of effort is being devoted to this **area** at research institutions and companies such as Netscape and Microsoft. As a result, many HTML enhancements have **appeared** such **as** *JavaScnpt,* DHTML, XML etc.

Although they improve the dynamic features of the web and make the web more interactive, none of them addressed the temporal demands of the web in a natural way. For example, XML provides developers with their own special purpose tags that can be defined for a variety of purposes. DHTML extends HTML by allowing multimedia applications to run on a user's desktop without interaction with the server. It also supports interactive databases and documents. However, the above enhancements don't provide a natural and efficient mechanism to allow servers to support complicated *remporal* phenomena in web applications.

With the introduction of Temporal Logic (TL), **THTML** provides an efficient solution for most temporal demands of a **web** site. It **has** great potential for adoption in complicated time-critical web applications such as the stock market, newspaper. web comrnercials etc.

THTML enhances the following features to HTML:

- **1.** THTML provides a mechanism to ailow one to define **arbitrary** venions for the same HTML file. **lt** aiso ailows one to define versions for parts of a page.
- **2.** THTML also ailows the author to **define** multiple instances of an HTML page, each of which **has** a time interval starnp. This time interval stamp determines what content should be posted based on the URL request and the request time from a client.
- 3. **A** tirne point can **be** specified by the client to indicate the specific instance of an HTML page it requires. The time **cm** be a *relative* point to now such as yesterday or the day before yesterday, or an *absolute* time point such as May 3rd, 1998 12:12:12, etc.

1.1.1 THTML Solution

THTML is a web-authoring tool intended for time sensitive web sites.

Time critical web sites have two basic requirements. One is that the page being viewed by the reader (the people who access or browse the site) is dependent on the time. For example, on a course web site, an instructor often wants assignment one posted during the first **week.** The solution of the assignment is posted in the second week dong with the problems for assigrment 2 and so on. The course notes should also **be** posted at appropriate times to match the schedule of teaching. In a newspaper site, different content should appear **every** day, and sometimes the reader would like to view yesterday's page as well as today's page.

The other requirement is that the content can **be** updated at customizable intervals, preferably automaticaily. This is often criticai for stock markets where a user would like to see the most recent content of the site.

THTML fulfills the two requirements with a solution based on discrete temporal logic. THTML allows a single source file to specify multiple *instances* of the corresponding page – different versions whose layout and content vary with time. Every request for a particular page in a THTML site is accompanied by a *requesf tirne,* which specifies the instance of the page that the reader wishes to view. The request time is often the **current** time but may **dso be** a time in the past or even the future.

The *author* can specify multiple instances for a *source file*, each instance associated with an *interval label*, in which the author can specify a time period or time point or sets of time periods or time points. When the server receives a request for a **URL** from a

Figure 1.1 Yesterday's Instance of a URL

reader, the *instantiation* (the process of taking the general definition of a time-varying

page, combining it with a particular time-point and producing the instance) of the source file begins. The page of the **URL** is configured dynamically with the THTML source file of each component whose time interval label is most relevant to the request time according to *the temporal versioning algonthms.* The page will **be** issued to the client when it is created, and discarded afterwards. The **dynamic** *configurarion* and the *temporal versioning algorithm* guarantee that the resulting page sent to the client is an instance at the time context specified in the reader's request time. Suppose, that a reader accesses a newspaper **URL** which has the following components: the title, banner, body and footer and requests yesterday's instance. The instance of each component corresponding to the reader's request time, namely, *the local* **time** *the server receives the request minus one day* in this case, constitutes a page for the requested URL and it will be issued to the reader as the response. An instance of each component is generated using the THTML source file that is most relevant to the request time.

Further, web pages change as time elapses, the instance a user is viewing may becorne obsolete. With push technology, synchronization *cm* **be** achieved between the client and the server. After the initial request for a **URL,** the server **sends** updated instances to the *reader* at frequencies that the *reader* specifies.

1.2 The Relationship between IHTML and THTML

The implementation of THTML is based on Intensional HTML 2.0, which supports sites having multiple user-defined versions. **THTML** 1.0 extends it by allowing users to create temporal variants of each arbitrary user-defined version and associate time with each temporal variant.

Because THTML adds a time dimension to the IHTML version space and specifies more time related version control algorithms to eliminate temporal version ambiguity, THTML makes the web site more dynamic and time-sensitive.

1.3 Terminology

A few terrns are used in particular ways in this thesis and the meanings are

explicitly declared here:

Readers-- the people who access/browse the web.

Authors-- the people who write the source for the web site.

Source *files*-- the source which the author writes and which generates pages and components.

Time point -- a complete temporal context, a particular point in time.

Instance-the version of a page/component/site corresponding to a particular time point.

Instantiation $-$ the process of taking the general definition of a time-varying page, combining **it** with a particular time-point and producing the instance.

Time interval stamp—the time periods or time points specified at the beginning of a THTML source file, which is the basis for determining the most relevant instance to the reader's requested time. It is distinguished with a tag pair $<$ cron $>$ and $<$ / cron $>$.

Temporal versioning algon'thm-the rules for determining the most relevant instance of a source file.

1.4 The Structure of the Thesis

The relevant background knowledge is introduced in Chapter 2, including the Intensional Logic based version control approach. Knowledge needed to implement THTML is also introduced, namely the *HTTP* protocol and the client/server model for communicating over the Internet. Chapter 3 examines THTML from three different perspectives: Language Design, System Design and Algorithm Design. Further enhancements to each design aspect are also discussed. Chapter 4 shows how to implement **THTML** on a web semer using the *Apache* **1.2.5** *APL* Chapter 5 concludes the thesis and proposes further directions for research.

Chapter 2

Background

2.1 Intensional Logic

In natural language, there are **many** situations where the semantics of a sentence depends on implicit contexts. For example, in the phrase

Today **'s** *temperature isfirpe degrees less than pesterday 's temperature,*

the value of today's temperature relies on the value of yesterday's temperature at the same place, therefore, the sentence *"today's temperature is 5 degrees*" may or may not be true.

Although implicit context sensitive expressions have for a long time been considered a non-mathematical and illogical aspect of naturai language, logicians have **tried** to capture them in formal methods. Intensional bgic is the result. It is **the** study of context sensitive properties using formal methods. Intensional Logic uses predicate calculus to represent these context sensitive expressions.

An Intension is an entity that depends on implicit contexts. For example, temperature is an intension that is variable in time and space. The time and place that is implicit in a sentence determines the temperature's value. In IL, *intension* and *extension are* distinguished for an expression. It **is** based on the distinction between *sense* and *denotation* of it. *Intension is the entire concept of an expression—what we, at some fevel, intend when we write it, while the extension denotes the particular object it in fact denotes (or currently denotes). For example, the expression " the President of France" denotes currently Jacques Chirac; but no one would claim that M Chirac somehow sums irp rlie rvhole concept of the French presidency[l8].* **As** *we* can see, an intensional object may have different extensions at different implicit contexts. For example, the French *consriturion specifies that the French President is directly elected; this is not the same as specifying that M Chirac be directly elected*[18].

The basic concept in IL is the *possible world semantics* proposed *by* Montague.

Simply speaking, it refers to a space of indices that an entity varies over -- the implicit contexts. **A** simple example is the World Wide Web **(WWW).** The web can be perceived as an indexed family of pages, the indices king the **URLs.** (URL's are required by the protocol to **be** unique, a necessary condition for indexing.)

Another example is the *chessboard* world, which is composed of al1 possible configurations of pieces on the board. As one can see, there may be many possible worlds in a given scenario. Using particular rules, one can switch from one possible world to another. (Modal logic is a variant of IL, in which people study the necessity and possibility of a statement in some or al1 the possible worlds, **and** the accessibility among them. Refer to **[l]** for more information in Modal Logic.)

2.2 Temporal Logic-An Instance of Intensional Logic

Temporal Logic (TL) is an instance of IL in which the collection of contexts models a collection of time points. Many areas in computer science adopt it extensively because high-level languages need a way of describing behaviors of complex systems in time. TL is useful for reasoning about a changing world. Since the temporal order of actions/events can be described, they are useful for representing dynamic systems. **TLs** are widely used in the specification and verification of reactive systems **and** in applications where the concept of time is central, such as temporal planning, temporai representation **and** temporal databases. Consequently, programming languages that provide access to such temporal concepts built on temporal **logic** such as *Chronolog* have a wide range of applications.

Further, the logic power of the language allows people to express complex temporal properties of the system. For example, if Discrete Temporal Logics (DTL's) are used, where time is represented as a sequence of distinct moments, a temporal formula can be used to effectively represent the individual steps of an execution. **THTML** adopts DTL so that the time context of a page can be efficiently reported at discrete moments.

2.3 An Application of Intensional Logic to Version Control

Plri Intensional Programmine Language **(PL),** which **can be** textual or visual,

retains two aspects of IL. First, at the syntactic levei, îhere are context-switching operators, cailed intensional operators, that can switch intensional contexts from one possible world into another. Second, at the semantic level, possible world semantics must be provided for each intensional operator. THTML is a web authoring tool that interprets version phenornena on the web using IL and provides an efficient method to generate instances of a **HTML** page and allow clients to switch arnong different possible worlds. THTML uses the **eduction** model, which is a tagged, demand-driven dataflow model. In this model, an intension is computed lazily. Each demand is tagged by an index. and is evaIuated by the extension at that index, the demand and their results flow as packets in an asynchronous network. in the eduction model in THTML, the **server** will generate specific instances and respond to clients only if the client requests it at a given time.

2.3.1 Existing Problerns in Version Control

Many version control systems exist that are very successful in solving particular problems. For example, pure version control systems, such as SSC and *RCS[4],* use delta techniques to solve storage problems and keep track of changes made by different programmers to a file. Software configuration systerns such as *MAKE[4]* aliow for automatic configuration of a system when changes are made to a component.

The integration of hierarchical-structured entities and version control is still not satisfactory. Using a system such as RCS and SCCS, for each file, there is a tree of revisions. The truck is considered to **be** the **main** version, the branches correspond to *iyariarrts.* Often, the changes are *merged* back into the trunk. The tree structure does not show how this merge took place. **if** integrated environments, such as **Adele[4],** are used, then for each module, there can **be** variants of the specification. For each specification, they can **be** variants of the implementation. And then for each implementation, there is an RCS-like structure for the development of implementation. In both cased, a tree structure is used for versions, yet, the tree structure is not appropriate for software development, because of the constant merging of different changes to the same system. A direct acyclic graph would **be** more appropriate. For example, suppose a program is

written to work with a standard screen in English. Two people independently modify the program. The first adds a graphics interface, and the other changes the error message to French. And then someone **asks** for a version that has both graphics and French messages. The new version inhents from its two ancestors. The main weakness of existing tools is that the different versions of a component have only a local signi ficance. It might **be** the case for example, that there is third version of component **A** and also a third version of component B. But there is no a *prori* reason to expect **my** relationship between the third versions of separate components. This lack of correspondence between versions of different components makes it difficult to build a complete system automatically.

An approach proposed by W. W. Wadge and J. **A.** Plaice based on IL provides an efficient rule for generating a version **[4].** In this approach, variant concept is addressed. It presents the need for versions of complete systems, and informally presents the versions of complete system can be generated by the versions of components. The approach introduced here will give the users the freedom of building any desired combination, but it is the user's responsibility of deciding which of the number of possible combinations will yield a consistent, working instance of the system.

The advanrage of *the approach is rhat it is now possible to talk of versions of rhe complete system—formed, in the simplest case, by uniformly choosing the corresponding versioris of the cornponents[4].* Suppose, that we have created a *fasr* version of every component of a compiler. Then we build the *fast* version of the compiler by combing the fast versions of all the components. Of course, in general it is unrealistic to require a distinct fast version of every component. It may be possible to speed up a compiler by altenng only a few components, and only these components will have *fat* versions. So *we* extend our configuration rule as follows: to build the *fasr* compiler we take the *fast* version of each component, if it exists; otherwise, we take the ordinary *vanilla* version. As generalized, in this approach, a system is perceived as a hierarchical structure that is composed of components or modules, each of which has its own set of versions. The system is configured with the most relevant versions of each component.

2.3.2 Intensional Solution

The following mechanisms are adopted to configure a particular version of a system.

2.3.2.1 the Intensional Configuration Rule and the Refinement Order

A partially ordered algebra is defined arnong versions, it reflects their refinement relations. $V \subseteq W$, *read as "V is refined by W", or "V is relevant to W", means that W is the result of further developing version V. The basic concept is that in configuration version W of a system, we can use version V of a particular component if the component* does **rlot** *exist i)~ a more relevant version. That is, we cari use version V* **of** *the comportenr as long as the component does not exist in version V', with* $V \subseteq V' \subseteq W[4]$ *. The major* advantage of defining a refinement relationship between versions for each module is that the creation of a system version **can** be automated. Note that the most generic version of each component *(vanilla)* must **be** created as a last resort in the system in case the system cannot find more refined versions of a component.

2.3.2.2 **Subversion**

If a piece of software is going to be adopted in different environments, it musr meet the different requirements of the users. The differences could exist at **many** levels. For example. at the level of user interfaces, some prefer text menus while others prefer the visual interface and mouse, while at the functionality level, some clients are provided with full access to all functionality while others only partial access. Different in implementation may also arise from one machine to another. The versions for machines X and Y may be identical, but differ with that for machine *S.*

In this approach, *subversion* (variant) is introduced. It is partially addressed in SCCS and RCS, with the introduction of branches, unfortunately, replying on the numeric strings to identify the branches becomes very unwieldy. We choose the **path** of naming the subversions. For example, a user can define a *bugFr* version of *2.3.4,* and a *release* version of *2.3.4,* a *bluesky* subversion of an *image* version or an *indigo* subversion of the *image* version. Subversions of a version are especially useful for parallel development.

2.3.2.3 Version Join

By joining compatible subversions according to the refinement **rule,** different system versions can be created. The "+" sign is used to identify the join operation.

The concept of *least upper bound* is addressed in the joining of compatible versions, which refers to the most relevant version of a component according to the refinement relationship specified.

Consider versiorl **VI** *and* **VZ,** *version* **VI+ V2 is** *the least upper bound of VI and V2, if and only if for all V such that VI* \subseteq *V and V2* \subseteq *V then VI + V2* \subseteq *V exists. In other words, VI +V2 is the least upper bound (the most relevant) because of the axiom* $(VI+V2) \subseteq (V+V) \equiv V[4].$

As above, we can create a new version for exarnple, *Japanese+ graphies+ infinite* X, if all of the versions of the modules are compatible.

2.3.2.4 Discussion of the Approach

Based on the refinement relationships, a system version can be generated automatically, and the number of possible system versions is greatly increased, because this approach dlows users to combine arbitrary versions for each component.

The intensional semantics of this version control approach is that the version universe of the system contains arbitrary user-defined possible worlds. **A** version of software can **be** perceived as variables in the universe. It is configured by the most relevant coordinate in each possible world. Using this approach, the version universe is expanded with each increase of user-defined dimensions. This approach **was** successfully implemented in **MTML,** which extends HTML **by** allowing arbitrary userdefined versions for a site.

2.4 Applications of the Version Control Approach to WWW----Intensional HTML

Based on the TCP/IP protocol implemented in the Internet layer, HTML and HTTP are used between the clients and the server.

On a web site, **HTTP** is used as the basis for building a **tme** client/server multimedia environment. HïML documents contain **links** to different data types. Web

browsers use plug-ins to interpret the different types of data and display it on the screen. This multi-media environment can consist of ail kinds of information that a cornputer can process.

HTML documents contain markup codes cailed tags in their body to control how the text is displayed. The markup tags tell the browser how the marked up text should **be** displayed—for example a header, link, bullet, list or body text. They can also contain other information such as a **URL** or the file name of an embedded picture. The format and usage of the different tags are detailed in the following sections.

2.4.1 IHTML Introduction

MTML introduces a version control mechanism to HTML by providing an efficient representation of the web's intensional characteristics. (See more information for IHTML in reference [3].)

In MTML, a **web** site is a hierarchical structure composed of many user-defined modules, each of which has multiple user-defined versions. **A** particular version of a web page is composed of the corresponding versions of each module.

In THTML, a version is specified with a *dimension name* and a *version value* in that dimension. For example: *bgc:blue* indicates a version is blue at the user-defined dimension: background. Join is represented **as** a "+". A version of a web site **cm be** *bgc:blue+ lang:english,* which is composed of the most relevant versions of *blue* in the *background* dimension and *english* in the *language* dimension. In all IHTML systems, a vanilla version, named **aai.html**, should be defined for each module.

IHTML follows the configuration rules introduced in the above section. IHTML is implemented by keeping only one copy of al1 variants; the server generates the final system version dynamically using the most appropriate version of each module. After sending the resulting HTML files to the client, the server discards it, so that maintenance of these versions can **be** avoided.

From the intensional semantic perspective, IHTML provides context-switching operator. This allows **users** to switch contexts and modify versions of each element on **web.**

2.4.1.1 Intensional Context Switches in Web Elements

A web element is an entity defined with a special tag and attributes in HTML such as a h yperlink, SSI. In HI'ML, because context switch operators are inserted into the web elements, web elements become intensional web elements, each of which has an intensiond and extensionai expression. The intensional context switches are implemented in hyperlinks, images and links in IHTML.

Intensional Context S witches in Hyperlinks

Extensional expressions in hyperlinks for example, have the following format:

 highlighted text ****

in which the absoiute version of the **URL** is given no matter what context the current **paze** is in. This means that, for example, if the current page is in **lang** : **f rench+pic :big,** then the version of the above hyperlink is modified to **lang** : **english** defined by the VERSION tag.

Intensional context switch in a hyperlink defines the version or dimension modifiers of the Iink so that the user cm switch to the relevant version based on the current one. In the above example, if the VERSION is changed to **VMOD,**

**<a href-nVRC-rd&asma VldOD=al~g:.nglimh~ tmxt **

the version of the **URL** therefore becomes **lang:english+pic:big,** which is **the** result of the merge of the cument context **lang** *t* **french+ pic :big** and the local context modifier on the *lang* dimension, namely **lang : english.**

Intensional Context Switches in Image Elements

Context switching operators are embedded in an image elernent in **the** sarne way as h yperlin **k.** The extensional and intensional expressions of an image are distinguished with **VMOD** and **VERSION** respectively. For example:

< **img ~rc-~iaugm.gif~ VXOD=mluag:frm~chaw Frmach varmion ,** which defines **the** version of the **image .gif** based on **the** current version of the page, while

 \langle img **src="image.gif" VERSION="lang:french">** French version ,

which defines the absolute version of the **image**.gif to lang: french.

• Intensional Context Switches in Server Side Includes (SSI)

MTMIL also **has** a Semer-Side *Inclrrdes* feature that causes the contents of the included file to be incorporated into HTML page **at** the server before it is sent back to the client.

The advantage of it is that the included information is included on the **fly** at request time.

A **SS1 may** have multiple versions, the files included in semer's response wili **be** the best-fit version for the user's request. For example

<!--#includm ~irturl=~hordmr.htmi~ -- >

if the page's version is currently French, the file included will **be french** version of

Figure 2.1 lHTML Working Model

header.html if it is in storage; otherwise, the vanilla version will **be** sent instead. (An error message will **be** sent to the user if the *vanilla* version of the target file cannot be found.)

2.3.1.2 Multiple Candidates and the Best-fit

The IHTML system keeps multiple versions of each web module. When the server receives an IHTML file from a client, the file pools are searched for the one matching the client's request. If more than one fits the request, the best-fit version will be found according to the ambiguity-handling algorithm. Figure 2.1 illustrates the IHTML working mechanism.

2.3.2 Surnmary of IHTML

As an IP language on WWW, IHTML integrates intensional semantics into the HTML and context-switching operators for the user to switch among different contexts. There are many advantages of IHTML over HTML.

Figure 2.2 Typical Client/Server Mode1

First, from functionality perspective, it provides a natural method to support arbitrary user-defined versions for a site, the user can switch versions for an element or a whole page by context switching operators, and the version of each **can be** totally or partially modified. JHTML provides great flexibility to the user in version

modifications.

Second, in the implementation, it saves system storage greatiy by keeping versions of variants rather than the system versions, in which there may be cloning parts among them. The server dynarnically creates a page by the configuration rule of the system version mentioned in Section 2.3. The version configuration of a site is automated to the greatest extent.

Third, it provides a base from implementation of THTML, which will **be** presented from the next chapter.

2.5 Base of THTML 1.0

The software of **THTML,** is implemented on APACHE server **APL** In order to introduce the implementation of the THTML system, **HTTP** protocols, which the HTML protocols and the servers' behavior are based are introduced first as the background knowledge.

2.5.1 HTTP Protocol

HTTP protocol is based on TCP/IP. The typical web model is client/server model, in which, the client initiates a TCP connection with the server, after the connection is established, the client sends a request down that channel. The server examines the request and responds in a manner specified in the server **Apache** *MI* and third-party modules. When a user types in a URL, the protocol name, (http in our case), server's *name(www.somewhere.com), the directory name(/where/foo.html)* and the *port number* **(use** *the defartlt ifnot specified,* **80)** are used to establish a unique connection with the server. The request for an **HTML** document is issued to the server on the dedicated network with the following format:

GET /where/foo.html **HTTP/1.O<CR><LF><LF>**

GET defines the method of retrieving a resource (norrnally a file) from **the** remote host. **/some/where/foo.html** defines the virtual host and the directory name, i.e. where the file is located. **HTTP/1.0** is the protocol name. The server sends the file back if it can find it on the server or decline the request if the resource does not exist. **If** no more requests happen during a certain time, it will close the connection to the client.

2.5.2 Server-Side Software--Apache *Server API-1.2.5* **[17]**

In THTML, special actions are defined at the server to handle THTML requests from the web clients. The software is implemented by extending a module based on Apache API server software, which offers a high degree of drop-in compatibility with the **popular** NCSA server **[2].** Internally, an APACHE server is built around an API (application programmer interface) chat allows third parties to add **new** server functionality. Most of the server's visible features (logging, authentication, access control, CGI, and so forth) are implemented **as** one or several modules. The implementation of request/response handling is discussed in the following sections.

2.5.2.1 Common Server Behavior

At the back-end of the server, after the server software has been compiled and got running, a daemon process will **be** running in the background, listening to any requests for connection. After one is detected, it forks a child process to handle it. Based on the data structure received from the client, it will determine which host and port number are the targets of the request. The virtual host then takes the path inserted in the request, and reads against its configuration to decide on the appropriate response.

2.5.2.2 Apache Solution Features

Apache makes the server running with default behaviors or customized behaviors defined in users' configurations and third party modules. There are two factors determining the server's request/response, nmely the *directives* in configuration files and the *exrended modules* at the server.

Directives

Directives are provided in the configuration files to aIlow a user to specify hislher preference on semer's behavior. For exarnple, a *port number(8080)* and the server *name(valdes.uvic.ca)* have to be defined as well as a *log file* for THTML for recording the behavior of the server. These features are all conveniently specified in corresponding directives in a configuration file.

Module Structure

Apache handles the requests and response by several steps (phases). They are:

- **1.** URI to filename translation
- *2.* Several phases involved with access control
- 3. Determining the **MIME** type of the requested entity
- 4. ActuaIIy sending data back to the client
- *5.* Logging the request.

The user can extend modules to handle any or all the phases. Some phases go through al1 the modules linked with **it** before going to the **next** phase such as logging the request and handling access control, while others may stop if one of the modules linked with the step successfully handles the request such as the step of URI to file translations, in which if one module handles the request and successfully translates the URI to a file, it **wiII** stop and go to the next step.

In the THTML implementation, a module is extended to handle the translation of a temporal URL and sends clients a response. The implementation of THTML is presented extensively in Chapter **5.**

Chapter 3

Time in THTML

3.1 Overview

There are three components that distinguish THTML from conventional **HTML** for both the reader **and** author of web sites.

1. Time sensitive tags are used to allow a reader's browser to specify a request time. This request time specifies which instance of a page is desired.

2. The author cm create multiple intemai-labelled source files for a single page or component. These interval-labelled files indicate when a particular page should **be** posted.

3 The times in these pages and requests can either be relative or absolute representation.

Some questions still remain, however. on how the above will **be** implemented. How will time be specified such that it reflects the temporal requests from readers and authors? There is also the question of how a request with different possible matches is resolved. These questions will be answered in the foilowing chapter where we will in troduce the existing temporal modeis and their advantages and disadvantages. We will also discuss a time-interval refinement algorithm that will help to resolve any ambiguities resulting from different possible matches.

3.2 Introduction to the Temporal Interval Starnp and the Time Sensitive Tag

Two new components were incorporated into **THTML** to support temporal requests. They are time sensitive tags (TST) and time interval stamps (TIS). TST's are used to specify which instance of a particular page should be sent to the client. For exampie, if we wanted an image of the weather outside at some specific time, we would specify the exact time in the TST and it would be the server's responsibility to match that tag to a

particular image.

On the server side, it uses the **TIS** associated with each HTML file to determine which instance should **be** sent to the client. The server is responsible for deterrnining the best-fit HTML page for each temporal request and sending the resuits back to the client. How the server determines the best-fit page and the design of both the TST's and TIS's will be discussed further in the following sections.

3.3 TIS Design

3.3.1 State-of-the-art Time Representations

Although time is continuous in nature, two common views have evolved: *conrinrioris* versus *discrete* time. Continuous time is considered to be isornorphic to *real numbers, whereas discrete time is normally isomorphic to <i>natural numbers*. The discrete interpretation of time has been adopted in many real-time systems because of its simplicity and the relative ease of implementation.

There are two time models widely adopted. One is the *linear time* model; the other is the *brunclring tirne model.* In the *linear* **time** *model,* time is considered to **be** totally ordered, i.e. if two distinct time points t and t' are given, either t is before t' or t' is before *^r*chronologically. However in the *branching time model,* multiple time lines are defined and a chronological ordering only exists between time points on the same time line. The branching time model is often adopted in hypotheticd reasoning, where time in the future cannot **be** determined, therefore, multiple time lines are assumed and researched hypothetically. For simplicity, we only consider the linear time model in this thesis. Therefore, time points are totally ordered on this single time line. Nevertheless, it is possible for web designers to use a branching time model, depending on the application.

3.3.2 Time-point and Tirne-interval Representations in THTML

The time model is not the only issue that needs to **be** resolved, there **has** dso been a lot of debate on the most appropriate representation of time. The two most common

representations are the rime *interval* and the **rime** *point.* An obvious choice is the time point. One reason for this choice is its simplicity. Time points are isomorphic to numbers and therefore **easy** to represent. The **other** reason is **that** the computational complexity of dealing with time points is less than that of time intervals.

The main drawback of the time point representation is that it is cumbersome for expressing the fact that an entity holds over a time period. As long as time is discrete, every time span can be represented by a finite set of time points. While this argument is theoretically correct, it is **often** not practical. For example, to specify that a THTML, file holds for a period of time on Web, the time period would consist of a list of **dt** vaiid time points. This representation is especially cumbersome when **the** granularity of time is fine-grained (e-g. seconds or even finer) and the average time a file holds is relatively long (say months or years). An obvious way out of this dilemma is to use time intervals in order to capture the duration of a temporal file. In addition to modeling the time span, the time interval can also **be** used to mode1 time points if the start **and** end points are the same.

The main drawback of the time interval representation is that the domain of the intervals is not closed with respective to the usual set-theoretic operations (union, difference and intersection). For example, the union of the two intervals may yield two intervals or just a single one.

> **[3-** *IO]* **~[15** *20]=/[3-10],[15,20]] [4,10] ~[6-17]={[4-I* **7]}**

The behavior of subtraction (set difference) is even worse:

$$
[4-10N8-20]=[[4-8]\n [4-10N6-8]=[[4-6],[8-10]\n [4-10N1-20]=[]
$$

Depending on the operands, subtraction yields zero, one or two intervals as **a** result.

This property makes it quite awkward and expensive to process time intervals on the computer. [14] Therefore, a question arises--namely whether we should use the *pointbased* model or the *interval-based* model in *THTML.*

Because it is common for people to specify time intervals over which files are to **be** posted on web, the tirne interval representation is needed to represent the valid times of files. In order to avoid the complexity it may induce, THTML adopts the representation of a *union* of sets of time intervals rather than *difference* in order to simplify the implernentation. Time intervals are perceived as the shorthand of time points within the starting and ending points inclusively. They are both adopted in the definition of **TIS.** In THTML. a reader must supply a complete time point to retrieve an instance. THTML doesn't support request *intervals,* therefore, in the TST specification, only time points are adopted. However, both time intervals and time points cm **be** defined in the timeinterval stamp.

3.3.3 Time Pattern Design in Time-Interval Stamp

3.3.3.1 Temporal Data Type in THTML

The data type of the TIS should have the capability of presenting both a time interval and a time point. Before discussing the design of the data type, we first clarify a terrninological problem. There are two terms widely used to represent time: *date* and *time.* Sometimes, the distinction between a date and time is made based on the chosen granularity. **A** date covers the granularity *year. monrh* and *day* while tirne covers every time unit below a day. Therefore, for exarnple **1985/3/L5** is called a *date* while **13:45** is **called** a time. Depending on the chosen granularity, it may **be** more appropnate to use one term over the other.

Nevertheless, THTML combines the two **as** the time type, which is *time* + *date.* There are two main reasons for adopting this data type. First, an acceptable granularity for representing the valid times of a file needs to be provided. Normally, on a computer, the granularity of files is in seconds. Although there are some arguments that the development of files should **be** granulated to below a second, the speed of web **transfers**

is insensitive to precision **below** a second. So the granularity in THïML is a second. This includes the modified time, the expired time, and the user-request time. This corresponds nicely with most version control software that also records time with precision of a second. The second reason is portability. The THTML time-keeping mechanism is consistent with that of all computer systems in use today—Unix, Windows, MAC and DOS. (whose time granularities are all up to seconds) So THTML is not tied to any specific architecture or operating system.

The time type is specified using **sir** units. The format is as foilows:

second **minute hour day** *month year [dow]*

The **valid** range of each unit is the same as that in the time-keeping mechanism of Unix as foilows:

> *second* **LOO-591** *minute* [00-59] *hour* [00-23] day [00-30/31/29/28] *month* [1-12] $year$ [0-] dorv **(days of rueek) [O-61**

Legend:

Among the seven units. **dow** (day of the week) is optional, it is provided to facilitate the temporal specification. There are two reasons that the unit of **dow** is added. First, in some cases, people have a habit of specifying a time pattern based on the day of week. Second, if there is not sufficient information given by the preceding time units to calculate the date, the server calculates the date from the data in **dow.** However, if there is a confiict between the **dow** and a specific date of the month, the value in *dow* is ignored.

Some users have a habit of specifying some of the time units while ignoring others, so, \bullet is specified to indicate arbitrary values in the valid range of a data unit.

3.3.3.2 Time Pattern Definition

There are five basic types that are supported in temporal logic: *during, after, before,* **point** and *recrirring.* Especially, the *recrirring* pattern facilitates the web authors to specify web postings that occur periodicaily.

The precise definition and their representations are as follows:

During, after/before and *recurring* identify time interval patterns; whereas *point* indicates a time point pattern on a single time line.

During is a bounded interval specification. It occupies only one segment on the time line. "-" is used to specify the interval between two points. *After/Before* are unbounded intervals. "<" or ">" preceding a time point is used to denote the time **period** *ufier* and *before* a particular time point inclusive. *Point* occupies only one point on the time line. *Recurring* describes multiple time-points or bounded segments that appear periodically on the time line. It is provided especially for sites that change contexts frequently, such as commercial sites, ", "is used to specify a recurring time pattern. So for example, the time-interval stamp

> **10 30 52 1 2 1998** ⁺

indicates a file should **be** posted *afier* time point **10:30:52** February **1 1998** while

< **10 30 52 1 2 1998** ⁺

indicates a file should **be** posted *before* time point **10:30:52** February 1 **1998.** If **a** web author intends to post a file on every Monday in January, February, Mach **1998,** he cm write the following time tag in the TIS:

 $* * 1.2.3 1998. 1$

If he determines to set the time period from **Jan** to March **1998** to post **a** file, he cm write in the TIS:

+ + + + **1-3 1998,**

One can combine *recurring* patterns with *interval* patterns in usage. Note that with the **Il-** " in a row, the user can specify the *srarting* and *ending* points of *an* interval on the time line. For example, time-interval stamp

+ **1-2 3-4 5-6 1998 l.2,3**

indicates a file should be posted every *Monday*, *Tuesday*, and *Wednesday* in the time

3.3.3.3 Discussion of Normalization Operation

A normalized time pattern is one that consists of only the union of a set of time intervals. The process of converting all patterns to this form is called *normalization*. It has great significance in judging tractability and is also an efficient method for matching temporal requests with time patterns

Since *diference* operations are not defined at the moment, normalization is not needed in **THTML.** However, *diflérence* operations will **be** included in the future, and so ail time pattems must **be** normalized so that the same matching aigorithm can be used without any modifications.

3.3.4 Conclusion of Time Pattern Design

Although there are many data models that can be used to represent time, we have chosen a mode1 that is able to support most comrnon temporal requests and is aiso easy to implement. But the question arises, is it the best representation for a temporal Ianguage? **1s** there a measure for evaluating the expressive power of a temporal Ianguage? If yes, is it possible to have a canonical expression independent of various time representations? There are many debates on this topic, so far, there is no answer in the literature.

Nevertheless, THTML can still be improved in the following aspects:

1. One further direction is incorporating more complex logîc into the system so that the expressive power of simple time pattems *cm* **be** enhanced. Intensional rules, including mathematical algorithms, **can be** adopted to specify more complex temporal requirements from the web authors.

2. Another direction is to incorporate a data **flow** language (temporal prograrnming language) that describes rules for temporal data flow. For example, every other day can be expressed naturally and efficiently in dataflow **programming** language with a definition of time:

n= k fby n+2d; k= \$?ri *May* **3 1998**
Precisely, an infinite temporal stream of dates, namely every other day, starting from May 3, 1998 is defined. The backend support for it includes algorithms for parsing and for temporal rules.

3.4 Best-fit Interval Selection in a TIS

The definition of an interval refinement relationship is the principal design requirement in THTML, since time plays a key role in retrieving documents and ambiguity may occur if the time point in request lies in multiple time patterns in a **TIS.**

A TIS that brackets the specified time point is a *condidate.* If *a* single candidate is found. it is used to generate a conventional HTML page. If no candidates are found, the absolute source file is used. **If** multiple candidates exist, the best-fit candidate is chosen. The rule of retrieving the best-fit pattern is determined by *extensional* and *intensional* rules. With *DTL* (Discrete Temporal Logic) applied in THTML, an interval can **be** perceived as a set of discrete time points on the time line, then the basic rule (extensional rule) to choose the best-fit candidate becomes *the smallest sirbset encapsulating a given time point is the best fit.*

Since not al1 candidates can **be** perceived as subsets of another, the extensionai rule does not eliminate al1 ambiguities between time interval candidates. For that reason, we also provide an intensional rule.

3.4.1 Extensional Rule

THTML specities that the request time in a **TST** is a tirne point while the TIS **can** be either time points or any combination of time intervais. When a time request is received, the server searches for the best-fit interval. The best-fit interval is defined to be the smallest interval encapsulating a given time point. For exarnple, the two time patterns A and *B* below are both candidates for a time point P, because P lies in the intervals of both **A** and B.

B refines **A** extensionally because the set of points in B is a subset of the set of the points in A.

Figure 3.1 Comparison between the Temporal Patterns of the Same Type

• The comparison of "after" and "before" patterns is also simple:

Figure 3.2 Comparison between %fier" Patterns

Obviously, in Figure 3.2, the set of points in B is the subset of that in **A** to the time criterion *P,* therefore, B is the better-fit candidate. As the same token, in the comparison between two *before* patterns, the one that is the subset of the other is the better-fit.

The absolute time interval refers to the unbounded interval from negative infinity to positive infinity. A conventional HTML file is perceived **as** a THTML file labeled with the absolute time interval. Using the extensional rule, it is easy to see the *absolute time interval* is the most generic one among all candidates; for it subsets all time patterns. So, the HTML file is the most generic source among al1 **THTML** file candidates. It is used as the best-fit when ambiguity happens.

3.4.2 Intensional Rule

The intensional mle applies when comparing **as** subsets is not enough. It is designed to minimize ambiguities that could result from the extensional rule. For example, in the situation in Figure **3.3,** an ambiguity exists because neither of them is a subset of the other.

Ambiguities exist in the comparison of the same time pattems as well **as** different time pattems using the extensional rule. Let us first discuss al1 possible combinations of patterns and then we will show how the intensional rule resolves these arnbiguities.

Figure 33 Comparison between Different Type of Temporal Patterns

3.4.2.1 Ambiguity Resolution among the Same type of Temporal Patterns

Since the time tags in **TIS** are unique, no arnbiguity cm exist if the **TIS** consists of just a time point.

When two infinite **recurring** patterns are compared, ambiguities may occur if one is not a subsec of the other. For example. in Figure 3.4, whether A or B shouid **be** chosen as the best-fit is hard to judge.

To eliminate the arnbiguity in this situation, the intensional rule specifies that if the recurring pattern contains equal sized segments, the candidate with smaller segment is the best-fit. The ambiguity occurs in other situations.

• In the comparison between two *during* patterns, extensional rule applies when one bounded interval is a subset of the other. However, arnbiguity occurs if one is not a subset of the other. For exampie, in the following scenario:

Figure 3.5 Comparison between bounded "Interval" Patterns

If A and B occupy the same length in the time line, ambiguity occurs with request time P1 or P2.

3.1.2.2 Ambiguity Resolution among Different Temporal Patterns

• The comparison of different time patterns is more complicated than for the same time patterns. For example, it is hard to tell if a *during* pattern candidate is more refined than a *recwring* pattem candidate, therefore, additional refinement rules are needed in the intensionai nile.

We **can** deduce that the *during* pattern has *a* smaller temporal gap than the *before/after* pattern, because the *before/after* pattern indicates a time period with an infinite end. This has a larger gap than the finite intervai in a *during* pattern, therefore, it is more genenc in definition than the *during* pattern.

The time *point* has smaller time gap than any other patterns, therefore, it is more refined definition than other patterns.

However, there are some situations where it is hard to find which pattern is more generic in relationship to a particular time point if two different patterns are given, say, *before* pattern versus *an* **afier** pattem and a *recurring* pattem versus a *during* pattern.

Ambiguity occurs in the first situation because the comparison cannot **be** solved mathematically between the two infinite segments; while in the second scenario, one is an infinite time pattem and the other is finite time pattern. An infinite pattem is not necessarily less refined than a finite pattem because there may be some situations for which one of the segments of the *recurring* pattern is closer than the *interval* pattern to the time point.

The *recurring* pattern describes *a* file to **be** posted on the web periodically. The intensional rule defines that the *recurring* pattern is the *least-refined* among al1 patterns.

One reason why the recurring pattem is the least refined is because it cm **be**

represented by a set of *during* patterns. The user **can** just use this method if he/she wants to obtain a higher priority.

The *during* pattern is defined to be more refined than the *before/after* pattern. If instead the *during* pattern were less refined, one could not use any other alternative time patterns to give *during* a higher priority. (Note that we do not allow an intervai to **be** denoted as time points. Although it is feasible in theory, it is not practical to define a large time intervai with time points.) There is an example to illustrate the above reasoning. If two time patterns are given, one is

>**10 10 10 2 3 4 1988** +.

the other is

+ + + **1,2,3 1988** +

If one intends to give a file, associated with the second pattern, a higher priority than a file associated with the first pattern in case that both cm **be** candidates to a request time, he/she can specify the second time pattern as a *during* pattern.

> $• • • • 1 1988 •$ \bullet \bullet \bullet \bullet 2 1988 \bullet + **3 1988**

As long as the frequency of the file is not very high, the above method is feasible.

Consequently, the refinement order in the intensionai **rule** is as follows. From the most generic to the least:

recurring before/after during point

Note that the above order is specified in terms of the time interval patterns adopted in THTML, it may not applicable to **other** temporal presentations.

3.5 Bat-fit Intewal Selection among Multiple TIS's

Because multiple time tags are ailowed in a single TIS, confusion may occur to the *best-fit* interval among the THTML files. In THTML 1.0, it is not the result of cornparison among ail patterns in ail THTML files at once, but among the best-fits of each file. Precisely, suppose that P1, P2, P3 are time patterns in THTML file T1 while *QI* and **QZ** are the time patterns in THTML file **72,** Suppose also that the candidate (relevent) labels are PI, *P2,* **Ql** and *Q2* are al1 candidates to a given request tirne *RI.* The process of retrieving the best-fit is as follows:

First, the best-fit among PI and **PZ** in **Tl** file and the best-fit among QI and **Q2** in T2 are resolved respectively. The best-fit of each is, say, P' and Q' respectively.

Second, P' and Q' are compared. The result is the best-fit to the client's request tirne **RI.** The THTML file that associates with it is used to generate the instance. The extensional and intensional rules are adopted in both steps.

The result is the same as that of comparing once. The reason for two comparisons twice rather than one is to simplify the implementation. The algorithm is similar to *brtbble soning.* the smaller subset of the first two will **be** kept to compare with the rest. The more refined in each comparison is always kept to compare with the next pattern. If ambiguity occurs, the most refined is still kept to compare with the rest. If no other patterns are more refined than it, the absolute TIS will **be** adopted, otherwise, the srnaIlest subset pattern will **be** adopted as the best-fit.

3.6 Conclusion

The design and reasoning behind the TIS and the rules to retrieve the best-fit intervals were covered in this section. In order to simplify operations involving the TIS, a linear time mode1 was adopted. The mode1 supports both a time interval and time point representation. This was done so that it would be consistent with the different ways users specify time. Since stating time in tems of a date and time of day is more complicated to handle, the two elements are combined together into a time data type, which granulates time in seconds.

This time data type foms the bais for specifying time requests **and** defining time intervals. Time points are matched to time intervals using the extensional rule. This rule is based on DTL, which the time interval is perceived as a set of time points on the time line. The time interval that is the smallest subset among those encapsulating the time point is the best-fit. This rule however still leaves some arnbiguities. Most of these arnbiguities can **be** resolved through the use of the intensional rules, which define a **refinement order for the different time pattems. The application of these two niles in one file or among best-fit pattems of multiple** THTML **files allows a server to determine which source file to instantiate given a URL request.**

Chapter 4

THTML System Design

4.1 Overview

The **THTML** system can **be** perceived **as** a typical three-tier architecture. It consists of the *clients' browser, file system and web server*. Enhancements are adopted in all three of these tiers. From the clients' point of view, two features (TIS and TST) are added to allow the client to specify a time element. The server supports this enhancement by implementing a tree structure and adopting a **naming** convention that maintains the hierarchical relationship between the modules and the different temporal versions of a module. The server determines which page to return to the reader **by** using a *tenipural versioning* and *time irregidatity handling algoritIm* to generate the most relevant instance for a particular **URL** request. The THTML system can be illustrated by the communication among the three components (See Figure 4.1).

Figure 4.1 Tbree-tier Mode1 of THThfL System

Upon receiving a **URL** request, the server decodes the address **and** extracts the request time. The server then searches the file system for the best-fit THTML source file

for that request time. It will then generate a conventional HTML file in response. In addition to the enhancements to the conventional web model, in which the server passively responds to the clients' requests, the server *push* (related to client pull) technology has been incorporated into the THTML server. This allows the THTML web server to actively feed clients with the most updated information at user-defined intervals.

In the following sections, we will discuss the issues related to implementing the above features. First the language design will **be** introduced in Section 4.2. In Section **4.3,** the backend system design is discussed, followed by the aigorithm design in Section 4.4. The chapter will conclude with possible future directions for each of the designs.

4.2 Language Design

4.2. l Design Requirement

In THTML, a web page is assembled using components indicated in a THTML page. These components are single-versioned entities in conventional HTML, but have multiple variants in **THTML.** Semantically, **THTML** ailows a conventional HTML page or its components to have multiple instances corresponding to the different time points. The generation of a **URL** instance at a specific time is, in the simplest case, a point-wise operation. It is generated by collating an instance of each component evaluated at the specified time.

The following figure shows the hierarchy in a THTML system. Each web page contains a list of components (images, sounds etc.) to **be** included in the page. Under each of these components, there are multiple instances corresponding to the different **time** intervals during which that component should **be** displayed.

Since THTML allows multiple instances for each **URL,** there must exist a **way** of determining which instance the user wishes to retrieve. This requires enhancements to the syntax and changes to the semantics of the HTML language. The syntax is enhanced by adding TIS's to the source files and allowing users to specify time points dong with their URL requests. On the semantic side, it involves interpreting temporal context

Figure 4.2 THTML Enbancement into HTML

switching **operators and helping to evaluate the** URL **at a particular time context and generating a response.**

4.2.2 THTML Language Enhancements to HTML

As mentioned in Chapter 3, a THTML file differs from an HTML file in two ways, namely, **the addition of TST's and TIS's. A THTML file has the following format:**

> <cran> *temporal-intervalsat terns* </cran> **<hm> <titla>** -... **</titlm> <body> THTML** *body* < **/body> </hm>**

The **TIS** is inserted at the beginning of a conventional HTML file with a special pair of tags **<croa>** and </cran>. The THTML body is sirnilar to a conventional HTML page except that it can contain optional TST's.

THTML body : : = *THTML cornponents*

THEML components : : = *(conventional HTML components)* 1 *(conventional HTML components having time sensitive attributes*)

THTML does not add **any** new tags in THTML body, however it modifies the existing tags by adding a new optional attribute. A **TMOD** attribute is added to the most common tags, found in HTML files, to dlow them to specify a temporal context switch. The TST's currently implemented are conventional hypertinks, images and SSI's. This attribute *cm* be added to other tags in the future. There is no Iimit to the number or sequence of TST's in a **THTML,** source file.

4.2.2.1 Time Sensitive Tag (TST)

Theoretically, any web component can be *temporalized* (made to **Vary** in time). For exarnple, temporal applets and images cm **be** created whose response to the reader is dependent on a time variable. In THTML **1.0.** three components were made temporal. They are *images, hyperlinks* and *SSl's.* The **BNF** format of the *TST's* for these *web* components is specified as follows:

The-sensitive-rag (TST) : : = **[<!-4nciude virtual="name.html"** *TMOD=temporal_context_switch >* ϵ src img="imgfile" $\Gamma \text{MOD} =$ *temporal_conte~t~switch* 1 **]*

The syntax of the *temporal-context-swirch* (TCS) variable defined in **BNF** format is as follows:

 TCS : $:=$ " " " *temporal_term* " " " ;

temporal_term ::= temporal_base $[($ *"+"* $]'$ *-"* $]$ *n* $($ *"Y"* \mid *"y"* $)$ $[$ $[$ *"+"* \mid *"-"* $]$ *n* (``M''''m'') | $(\text{``+'''}\text{ ``-''})$ n $(\text{``D'''}\text{ ``d''})[\text{``+'''}\text{ ``-'''}]\text{``H'''}\text{ ``h''})$ $[(\mathbf{a}^{\mathbf{b}} + \mathbf{b}^{\mathbf{b}} + \mathbf{b}^{\math$ **(bbw916&w99)** $n:=$ *digits*;

 $temporal_base ::= \varepsilon | natural_l | absolute_l;$ *natural_l* ::= "**tomorrow"** | "yesterday" | "now"; *absolute-t* : := **"\$"** *Week-day Mon Date Hour* **":** " *Min* ":" *Sec* ; *Week_day* ::= "Mon"| "Tue"| "Wed"| "Thur"| "Fri"| "Sat"| "Sun": *Mon* ::= "**Jan**"| "Feb"| "Mar"| "Apr"| "May"| "Jun"| "Jul"| "Aug"| "Sep"| "Oct"| "Nov"| "Dec"; *Date::* = digits from 0 to 31; *Hour:: =digits* **from** *O to 23; Min::=digits from O to 59; Sec:: =digits* **fiom** *O to 59;*

Legend:

the time unit identifiers: **Y, M.** D, **8, 1, S** (capitalized or lower case) indicate temporal units: year, month, day, hour, minute and second, respectively.

As shown above, simple temporal terminology can also used in context switch attnbutes. The **terms yesterday, tomorror** and **now** are supponed in **the** curent version and their semantics will **be** presented in **4.2.3.**

4.2.2.2 Time Interval Stamp (TIS)

The TIS is the other feature added in the conventionai HTML. TIS's are distinguished from other components by a pair of tags **<cron>** and **</cron>**. The user can specify a specific time point, a time period or a set of time points and **periods.** Al1 TIS's can be specified with six time units as follows:

second minute **hour day** *month year* **[dow]**

Each unit is separated with white space. The days of **the** week **(dow)** is optional. Each unit can be specified using a combination of **numbers** and symbols. The following symbols **can** also **be** used:

- denotes a time interval-
- , separates a list of valid inputs.

denotes any arbitrary value in the valid range of a data unit.

Therefore, for example, the time interval from May 3, 1998 to May 20, 1998 can **be** represented as follows:

<CZOD> + **3-20 5 1998 c/cronw**

1.2.2.3 Creation of a THTML file

THTML 1.0 provides a command in the UNIX environment for an author to create and modify a THTML source file. The command ivi originates from IHTML 2.0 and was used to create an MTML file **(See** [IO] for **more** details.) The author enhanced the command to create a THTML file. The format is as follows:

ivi $[-t]$ $[-\text{c}$ c c c d \text

The command ailows the user to edit the specified THTAML source file. When it is invoked, it will open the default editor. (specified by the **\$EDITOR** environment variable) The -t option is used to create a new version of *filename*. **html** using a copy of the rnost recently created version. This new version will be saved in the fil *ename.* **html** directory in the following format: **aai** . n. **htmi.** where n is a unique number generated automatically by the system for each version. An existing version can also **be** edited by specifying the version number in place of *cnumber>.*

4.2.3 Language Semantics

Formally, if each conventional HTML page is P and the request time is u , then each THTML page can be specified by (P, u) . At request time u, a unique instance of P is issued to the client. From the semantic point of view, the response to a URL request is an HTML page that is dependent on the time point specified. In other words, the URL is an intension. The response to the reader cornesponds to the process of generating **an** extension from the intension. An extension to a particular intension, namely an instance is generated using two components. One is the URL or HTML component, the other is the request time. The temporal contexts are not necessarily the client's request time. The

reader can define **any** time point. The web elements in the response for the **URL** inherit the context by default. However, they can also **be** modified locally by temporal context switches.

4.2.3.1 Semantic Distinction between HTML and THTML

4.2.3.1.1 The Global Context and the Local Context

Figure 1.3 Server lnterpretation Mode1

The *Global Conte-rt* and *Local Contexz* are distinguished by their scope. Because the context for the **URL** is adopted as default to interpret those for the web elements for the HTML page. In other **words,** the elements in the HTML page inhent the context of the URL by default. However, a context switch operator **cm be** defined in each local element to modify the global context. It changes the six time units as necessary.

At the server, a THTML document is converted into conventional HTML. Each component is instantiated using its local context. Figure **4.4** illustrates the conversion process.

4.2.3.1.2 Temporal Context-S witches *(TCS)*

Global contexts are converted to local ones by *temporal context nvitches.* These temporal context switches are implemented only in *hyperlinks, images* and **SSl's** in THTML 1.0. The usage and semantics of them are covered as foilows:

Temporal Context Switches in Hyperlinks

A temporal context switch **can be** specified using the attribute **TMOD** in a hyperlink:

When the hyperlink is invoked, a conventional **HTML** file, for the URL, is instantiated using the context specified in TCS. The locd context of a link containing a

Figure 4.4 the Local and Global Context

TCS is a combination of the base temporal context and the **local** temporal modifier. If a base temporal context is not specified, the global context of the curent page is used.

The temporal words **tomorrow, yesterday** and **now** are implemented in THTML as following manner:

Tomorrow is equivalent to the local modifier "+ ld". In other words, it represents the same time (the same hour, minute and second) as the default context, but increments the day by 1.

- **Yesterday** however has the opposite effect and decrements the day by 1.
- The keyword **now** however denotes the **default** global **context.**

These three temporal terms are interpreted in the same manner in the other eIements.

Extensional expressions cm also **be** specified in TCS's. For example, **<a** href="URL_Address" TMOD="\$Tue May 19 12:13:14+2H-1D"> denotes that the local context is changed to May 18 **10:13:14.**

Temporal Context Switches in Images

In an image element, the TCS's can change only the contexts of the included image entity. In other words, they perform no modifications to other web components. For example,

<- **mrc- mgraphic~-filem TXOD=mtwrrowm>**

If the page is invoked using a hyperlink that has the context modifier **yesterday** (ie. *the client's request time* $-ld$ *)* and the image link has the local modifier **tomorrow**. the image corresponding to the client's request time will **be** instantiated. (ie. *client's request time -Iday+Iday)*

Temporal Context Switches in Server Side Inctudes (SSI)

As mentioned in Chapter 2, SSI is a feature that the **APACHE** server supports. It alIows the author to include an HTML file in another **HTML** file, instead of copying the same contents. It facilitates web content creation and heips with version control by avoiding redundant code scattered among separate files. Temporai SSI is similar but also allows multiple variants in the time dimension. The TCS is inserted into a SSI component and modifies the local context of the entity. **It has** the following format < **I** --

#include virtual r -included-filen TMOD = + **offset>**

The *offset* can be specified using the *six* time units: *year, month, date, hour, minute and second* and is relative to the default temporal context of the *TCS*. (See 3.2.1 for The BNF format.)

A local temporal modifier and an absolute time point constitute an extensional representation of a SSI with TCS. Given <!-- #include virtual="file.html" **TMOD="** $$July 5, 98 12:00:00+7d +4s" >$, an instance of "**file.html**" at the time context: *July* **12, 98** *12:OO:W* should **be** included in the HTML page generated.

1.2.3.2 Web Intensions and Web Extensions

A web intension is the relative expression of a web context in terms of a default context. It is specified in a temporal web element with a temporal modifier; while a web *extension* is an absolute temporal context value. It is represented using an absolute time and an optional temporal modifier.

As introduced in Chapter 2, the sarne intensional expression with the sarne literal expression can infer different extensional values depending on the context. Consequently, the same temporal **URL** requests can **be** resolved into different instances (conventional **HTML** pages) depending on the time context,

4.2.3.2.1 "Now" Discussion--1ntension and Extension

Strictly speaking, there are 24 extensional representations for *now* because of the 24 time zones in the world. For example, given the following **TST:**

Example 1 Follow Constraint Constraint Analytic State of the Client lives in a place where "now" is *July 13*, 1998, 3:30:00 and he/she issues a THTML request to a server in another time zone, whose time is *"Jrtly* 13, 1998, 6:30:00", which time should **be** chosen as the default temporal context?

In order to deduce the right answer, let us examine the request process at the server. The user-request time is sent to the server in **GMT** (Greenwich **Mean** Tirne) format according to the request protoçoi. Since no mechanism is provided to retrieve the reader's time zone, then the local representation at the server becomes the default temporal context.

Therefore, the time related process at the server side is as follows:

When the server receives a THTML file, the reader's request time is computed and converted into the server's local time (which forms the basis for retrieving the best-fit version). The instantiation process begins after the global context is resolved. In the previous exarnple, the default time context is therefore *July* **13, 1998** *6:30:00+ld-211 =Jdy 12,* **1998 4:30:00** .

If a reader expects to retrieve an HTML page at the server with the same

extensional time representation as the client, the extensional representation of now at the client's side must **be** wntten into the temporal request. For this case, it should request

TMOD="\$Fri July 12, 1998 3:30:00 -1d+2h"

The process of conveying temporal information between the web client and the server is illustrated with the following graph:

Figure 1.5 Time Conveyance in THTML System

4.2.1 Conciusion

With the incorporation of IL, the server's functionality can be perceived as converting an intension (the **URL** or THTML tag) into its extensional representation (the web page generated in response). Temporaiized web elements are intensions that use TCS's to switch among different possible worlds. Web authors are allowed to define multiple source files for each elernent. Their TIS's detennine if it is adopted **as** the extensional representation of the elernent at a particular time context. Due to the above designs in syntax and semantics, the range of **HTML** pages viewable at the same **URL is** increased dramatically from the client's perspective. 4.3 Backend System Design

4.3.1 Design Requirements

One of the design requirements **was** to make the THTML language independent of web browsers. The browser should not know whether or not the server is a THTML or regular HTML server. So the temporal enhancements to the language rnust **be** supported only through the backend. In order to support THTML, the server must implement the following:

The system should have a mechanism for maintaining multiple THTML source

files for each web component.

It should dso maintain the hierarchical relationship between al1 THïML sources and their respective components.

A time tag is associated with each TIFTML file, so that the server **can** determine which file to retrieve based on the time context.

The above requirements are met in THTML *by* implementing a *rree* structure to maintain the hierarchical relationship, adopting a narning convention to differentiate between multiple source files and including **TIS's** in each file. Each of these aspects will be discussed further in the following sections.

4.3.1.1 Tree Structure

First, in file system design, a *rree* structure is used to maintain the hierarchical relationship between THTML source files and their respective web components. **A** normal URL indicates a directory in the **THTML** system that holds multiple **THTML** sources for that **URL.**

There are other mechanisms to maintain the hierarchical relationship in version space besides the *tree* structure. One commonly used method is to keep the **file** that contains the relationship between the THTML sources and web components. There are several reasons why it was not adopted. First, the web server wouId have to access the file for each instantiation process. The file would become a key component of the system and would requiring special maintenance. If it became corrupted, the whole system would shutdown. Furtherrnore, if a search for a target THTML source file was executed at the same time that the author **was** creating a new THTML fiIe, the file would become a *critical section* and a *synchronization* mechanism would need to be adopted.

The tree structure has significant advantages in comparison:

- By embedding the hierarchy into the structure it eases system maintenance
- The algorithm to retrieve the best-fit THTML source is simplified as well.

4.3.1.2 Naming Convention of THTML files

Second, a simple naming convention needs to **be** adopted to maintain the parentchiid relationship between a temporal object and each THTML file. A straightforward naming convention is described below.

When a file is created, it is put in its module's directory. For exarnple, **al1** THTML sources for the web component **home**.html are put under directory **home**.html. The format of a THTML source file is encoded filename.n.html, where n is a number generated automatically by the system when the file is created. It reflects the order of creation **and** distinguishes the different THTML files. An HTML instance, which is an extensional expression of a web component, has a file name ending with the suffix .html. It is similar to a conventional HTML file, but contains an encoded temporal context that is inserted between the file name and the suffix. For example, **home chodiaywbySqzGuTi_Rgej_nPkeai.html** is an instance of the module **home.html** at the relative temporal context: **osoxoa-iooiaorow.** The format of **an** instance is *rodule name*@encoded context.html. The encoding scheme helps the algorithm retrieve the best-fit THTML file.

1.3.1.3 Design of Time Interval Stamp (the Time Tag)

Third, the TIS **is** used to relate time to a specific THTML file. It indicates to the system which file is most relevant to a temporal request.

There are many reasons for keeping a time-interval stamp in a **THTML** file.

In order to instantiate an HTML page, the systern needs to keep time information for each THTML source file. Keeping the information in the file provides a straightforward way to link the time and source page.

It reduces the time to retrieve the corresponding time interval of a THTML source file.

It provides a naturai way to combine THTML files with regular HTML files. HTML files are just THTML files with time intervals from *negative infinity* to *positive* irzfiziry, which is denoted with time label + * *.

The TIS is the key to temporal design. It not only specifies the time relevancy of each file, but also reflects the refinement relationship between the different THTML fiIes.

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1.3.1.1 Conclusion of File System Design

The result of combining the tree structure, naming convention and time interval starnp is shown in Figure 4.6. Each **URL** corresponds to a module and therefore has its own directory. Each directory contains al1 user-defined and temporal source files associated with the module. They are stored using the naming convention <module_name>.<encoded_userdefined_version>. n.html. Although it is currently irnplemented on one cornputer, the design can be easily incorporated into a distributed file system.

Figure 4.6 System Structure Design

3.1 Algorithm Design

4.41 THTML Web Server Model

In a typical client/server model. when the client selects a temporal elernent on a THTML site, a connection is created between the client and server. The client **then** issues a request for an **HTML** file with a particular time point to the server and the server responds with an instance at that time point to the client.

In this model, the web server can be perceived as a black box. With the request of a temporal **HTML** file as input; the server outputs its extensionai representation.

During this process, a time irregularity-handling algorithm is used to obtain the correct local context from both the temporal request and the local **modifiers. A** temporal

versioning algorithm is then used to retrieve the best-fit THTML file. These algorithms will be described in the following section dong with the incorporation of *push* technology to keep the clients **updated** with the most current HTML **pases.**

Temporal HTML page

Instance Generated

Figure 4.8 Interpretation from THTML to Conventional HTML

4.4.2 Time Irregularity, User-defined Temporal Terminology and Hierarchy of User-Defined Calendars

4.4.2.1 Gregorian Calendar and its Irregularity

Calculations related to time have an irregular feature. For example, $\hat{\boldsymbol{\xi}}$ **Fri July 31, 1998** \div **1d**, should be *August 1, 1998* rather than *July 32, 1998*. The time irregularity is caused by the Gregorian calendar people use as an international standard for civil use.

A calendar is a **system** used to **organize** units of **time** for the purpose of reckoning time over extended periods. By convention, the day is the **srnallest** calendarical unit of time; the measurement of fractions of a day is classified as timekeeping. The generaiity of this definition is due to the diversity of methods that have been issued in the development of calendars.

In the Gregorian calendar, years **are** counted from the initial **epoch** as defined by

Dionysius Exiguus, and are divided into two classes: common years and leap years. A common year is 365 days in length; a leap year is 366 days, with an intercalary day, designated February 29, preceding March 1. Leap years are determined according to the following rule:

Every year that is exactly divisible by **4** is a leap year, except for years that are exactly divisible by **100;** these centurial years are leap years only if they are exactly divisible by **400.**

As a result, the year **2000** is a leap year, whereas 1900 and 2100 are not leap years. These rules can be applied to times prior to the Gregorian reform to create a proleptic Gregorian calendar. In this case, year O **(1** B.C.) is considered to **be** exactly divisible by *3,* 100, and 400; hence it is a leap year.

The Gregorian cdendar is thus based on a cycle of 400 years, which comprises 146097 days. Since the number 146097 is evenly divisible by 7, the Gregorian civil calendar exactly repeats after 400 years. Dividing **146097** by 400 yields an average length of 365.3425 days **per** calendar year, which is a close approximation to the length of the tropical year. Comparison with Equation 1.1-1 reveals that the Gregorian calendar accumulates an error of one day in about 2500 years. Although various adjustments to the leap-year system have been proposed, none has been instituted.

Within **each** year, dates are specified according to the count of days from the beginning of the month. The order of months and number of days per month were adopted from the Julian calendar **as** follows:

January 31 February 28 March 31 April 30 May 31 June 30

Jtily 31 Arigrtst 31 September 30 October 31 Novernber 30 December 31 In a *leap* year, February has 29 days.

Based on the representation of time keeping and the Gregorian Calendar currently used internationally, the defauIt hierarchical concept of the time units used in **THTML** is defined as *second* \leq minute \leq *hour* \leq *day* \leq *month* \leq *year.* A set of functions supporting the above concept is provided to guarantee the correctness of time calculations, which are discussed in next chapter.

4.4.2.2 User-defined Temporal Terminology and Hierarchy of User-Defined Calendars (not implemented yet)

User-Defined Temporal Terminology

As the design in further implementation, the TCS can be enriched with more complex temporal language with enhancement of additional backend supports.

For example, additional religion dates say *Easter* can be recognized in a TCS by incorporating the following operations to calculate the date of Easter in Gregorian calendar in THTML.

Y stands for Gregorian **year. Al1** variables are integers and the remainders of **al1** divisions are **dropped.** The finai date is given by two vaiues: M (the month) and D (the day of the month).

$$
C = Y/100,
$$

\n
$$
N = Y - 19*(Y/19),
$$

\n
$$
K = (C - 17)/25,
$$

\n
$$
I = C - C/4 - (C - K)/3 + 19*N + 15,
$$

\n
$$
I = I - 30*(U30),
$$

\n
$$
I = I - (U28)* (1 - (U28)* (29/(I + 1)) * ((21 - N)/11)),
$$

\n
$$
J = Y + Y/4 + I + 2 - C + C/4,
$$

\n
$$
J = J - 7*(J/7),
$$

\n
$$
L = I - J,
$$

\n
$$
M = 3 + (L + 40)/44,
$$

\n
$$
D = L + 28 - 3I*(M/4).
$$

By the same token, the user-defined caiendars and the temporal concepts can **be** specified with corresponding backend support. For example, if a student **can** use the term *semester* and **acadernic** year in his TCS, operations of parsing these terms and converting them into Gregorian caiendars should **be** provided.

User-defined Calendars and their Hierarchy

Funher, if multiple user-defined calendars are specified in the system, ordering should also be specified **among them** to eliminate ambiguity. For **example,** New **Yenr's Da?** has different Gregorian values in the Lunar Chinese Calendar and the Jewish Calendar. If a temporal request having a term *New Year's Day* is received, ambiguity must be resolved if both calendars **are** specified in the system. Therefore, an order should be specified between **them.** In the future, an intensiond tier for user-defined calendars, their ordenng and operations should **be** provided by the system.

The relationship of components in the intensional tier is as follows:

Figure 4.9 User-defined Caleadars, Orders and Operations

Figure 4.10 Intensional Tier Supporthg User-defmed Time Concepts

The relationship of the intensional tier with other components in the system is illustrated in Figure 4.10. The server will retrieve the temporal concepts from the intensional tier and compute the results.

1.4.3 Temporal Versioning Algorithm

The purpose of instantiation is to generate an **HTML** instance in response to the URL request. After the time context is **extracted,** the server will search for al1 **THTML** source files that match the time point specified. If there is only one match, that file will be used to generate the HTML instance. Otherwise a best-fit file is chosen using the intensional and extensional rules described in Chapter 3. If the rules cannot resolve all the ambiguities, then the absolute (default) THTML source file is used. This file is just an ordinary HTML file with no time stamp. **It** is kept in the system as a last resort when no time intervals match a specific time point.

The THMTL source file is then parsed. Any links, images and other tags with temporal attnbutes are also instantiated recursively until **al1** temporal elements are converted into their respective extensional expressions,

4.4.4 Push Technology and Dataflow Model of WWW

With the rapid growth of time sensitive information on the web, the reader needs to be periodicalIy updated with the most up to date information. This is facilitated in THTML with the incorporation of *push* technology.

4.1.3.1 Problern of Cache in the Existing Web Model

The cache is one of the most significant features of the major browsers (Internet Explorer and Netscape Communicator). With a cache. copies of responses are stored in the local file system and can **be** reused for the same **URLs** without successive requests to the server. Therefore, caching conserves bandwidth and reduces network latency. However, if the caching mechanism not implemented appropriately, there is a **risk** of receiving stale data from the cache.

In order to enhance cache control, the *HTTP* protocol contains both explicit specifications as well as heuristic hints in **HTTP/1.1,** conceming the cacheability of documents and how often their freshness should be venfied in order to guarantee that they are still up-to-date.

One of the most important features in **HTTP/l.O** about caching is the conditional GET. It allows a document to **be** retrieved conditionally, based on whether it has been modified since the last access. If the document has not been changed, a very short **not modif ied** message is issued; otherwise, the updated document is transferred. For example, after a client requests a document **/some/where/foo.html**, the following response can be sent to the client:

> HTTP/1.0 200 OK Server: Netscape-Enterprise/2.0 **Date: Sat, 19 wr 1997 IO :22 :O0 Oçr** Last Modified: Fri, 18Apr 1997 15:12:05 GMT Content-type: text/html **Content-length: 6510**

The above headers indicate the server and the version of HTTP protocols. The time the web transaction happas, the last modified date of the file. In the subsequent requests, a conditionai GET request now adopts the timestamp from the **Last**modified header and issues it along in the request header with an If Modified **Since** header to the server. The example of the subsequent requests is as follows:

> **Get /somm/whare/foo.html ETTP/1.0 If Modified Since: Pri, 18Apr 1997 15:12:05 GMT Accept: taxt/plain, taxt/html**

It indicates that if the document is not modified since the date defined in the **ff Modif ied Since** header at the server side, the document content is not issued to the client. Therefore, the following headers but contents will **be** sent the client as the response.

```
HTTP/1.0 304 Not Modified
Server: Netscape-Enterprise/2.0
 Date: Sun, 20 Apr 1997 15:45:12 GMT
```
However, if the document has changed. the server will feed the client with an updated document and a 200 OK response. 304 Not Modified response saves bandwidth and reduces latency, as no document transfer actuaily occurs.

The above mechanism is efficient in conventional **HTML** to keep the most updated content in a client's browser; however, in **THTML** file system, as the **time** changes, it is

usualiy a separate file (i.e. a different version) that should **be** issued to a client because the original document is not changed and stored as a histoncal version in the system. Therefore, the header indicating whether the file has been modified cannot prevent clients from viewing obsolete versions in **THTML** system.

3.4.4.2 Possible Solutions

Based on the above problems, two solutions are provided. The advantages and disadvantages of them are discussed in the folfowing sections.

Cache Control

The first approach is to delete the cache. **HTTP/l.** 1 allows a client to specify when the file expires and whether he/she uses cache or not. A typical adoption is web commercials. Normally, **THTML** documents do not change while the images within inline ads are. Therefore, the HTML content is not modified on a per-request or per-user basis in order to contain a different ad. Instead, only the inline images update with a particular frequency. With this approach HTML pages can **be** cached. Only the images themselves force successive requests to the server instead of being served from the cache. Although the contents **can be** updated with this mechanism, the disadvantage of this approach is that substantial web traffic is induced. Therefore, it is not feasible for it to adopt in a large, *dynamic* HTML file applications.

Server Push

The other approach is to complement existing web protocols with a new technology----*Push* technology (compared to *Pull* technology), which is considered a more promising solution.

The original pull technology (client pull) refers to a model in which a client requests, then the server responds while *push* (server push) refers to a system in which a client automatically receives information or applications from a network server. With **server** push, clients do not have to issue successive requests to the server for the most updated versions.

In a server push, the connection between the server and a client does not end until

the client closes the session. The server periodicdly feeds the client with the most updated version. Because multiple connections need to **be** kept and the responses to the previous connected clients should **be** issued periodically, server push burdens the server. However, the approach is preferred in THTML. The advantage of it is that web traffic is reduced to dmost hdf (theoretically for lack of successive requests from the client). It **was** shown in a year-long study involving **4,000** users in six Fortune **1Oûû** companies that just Iess than 20% of network traffic was raised with the use of push technologies while 100% with the use of the client pull technology. Another significant reason that *server push* is preferred is that no particular software is required with the client, because that just less than 20% of network traffic was raised with the use of public 100% with the use of the client pull technology. Another significance push is preferred is that no particular software is required with the funct

Figure 4.1 1 Server Push

With this technology, synchronization can be achieved between a client and a server with a pre-defined frequency. It not only avoids the fikelihood that a client may **be** viewing the stde information but aIso reduces web traffic tremendously. Although today *push technology* is still in its infancy, many commercial push products have been available within industry, and web technology is taken into a promising new research direction. With **rime** addressed, an ideal solution is provided with push in order to bring the increasingly *dynamic* feature to the World Wide Web.

4.4.5. Conclusion of Algorithm Design

For convenience, THTML dlows the author to specify time using **the** Gregorian calendar. This however creates a problem when adding or subtracting time due to the irregular way in which time is denominated in this system. We have therefore implemented a time irregularity-handling algorithm to perform these operations correctly. This algorithm can **be** extended in the future to handle different calendars and

temporal terms.

Once the time point **has** ken determined, the best-fit time interval must **be** found. The temporal versioning algorithm uses both the intensional and extensional rules to determine this best-fit intervai and resolve any ambiguities.

Push technology **was** also incorporated into THTML to facilitate the continual update of a reader's page. It involves executing the above calculations at user-defined intervals. Although this results in a heavier load on the server and other solutions may exist, it is the most efficient solution for some applications where up to date information is essential.

Figure 1.12 Backend Components of THTML

1.5 Combination with IHTML

As a side, **THTML1.0** is compatible with MTML. **THTML** files of a user-defined version can be created and modified. This indicates a user-defined version cm be varied with time in semantics. Clients can request a particular instance of user-defined version of a web component at a particular time point. In a THTML file having user-defined version, the TST's have the following formats:

```
TST: := \leq:--#include virtual="included_file" TMOD="TCS"
      VMOD="intensional context_switches" > <src img="imgfile"
           \textbf{TMOD} = \textbf{``}TCS" \textbf{VMOD} = \textbf{``intensional}\_context\_switches" > |<a hre f - a URL_Address" TMOD-a TCS " 
           VMOD=aintensional-context-swi tches" >] *
```
• The format of *intensional_context_switch* is detailed in [10]. The TCS has the sarne format as introduced in Section **4.2.2.1.** Therefore, for example, in a hyperlink element, a client can write,

.

in THTMLl.0, temporal instances are implemented as special **subversions** of a user-defined version, so when an element having **both TMOD** and **VMOD** attributes is encountered, the user-defined version is first resolved, and then the process of instantiation, namely, the instance of the user-defined version is resolved in the time dimension. The resulting instance is sent to the client*. In the above example, attribute **VMOD** specifies the local, user-defined modifiers: **pic :big** and **bgc tblue** of the global user-defined version of the page. If the global, user-defined context of the page **is lang** : **english,** the local user-defined context of the eiement will **be** pic **:big** ⁺ **bgc** : **blue** + **lang** : **english.** After the best-fit user-defined version at the above context **is** resolved, the instance of it will **be** resolved. Attribute **TMOD** modifies the temporal context into *current time-Id* for the instance. The **THTML** file whose tirne interval stamp best fits the temporal context is instantiated. It replaces the intensional element into its extensional expression.

In the **THTML1.0,** multiple HTML files of a user-defined version *cm* be created and modified in the system. The command for creating a new THTML file from a userdefined version has the following format:

ivi [-v version [-b baseversion] 1 [-tl *module* It indicates that the author cm specify a THTML file of a user-defined version of the module module. The option $\begin{bmatrix} -v \\ v \end{bmatrix}$ version $\begin{bmatrix} -b \\ b \end{bmatrix}$ baseversion]] is used in the sarne manner as in IHTML 2.0 to define a user-defined version. If the **version** version does not exist and the base version is supplied, it copies the **baseversion** version of the page to the **version** version.(see the manual of MTML. 2 [lO].) A THTML file of the user-defined version will **be** created if there is option **-t;** otherwise a

--

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^aThe semantic of versioning resolution principle is omitted. because the version resolution principle will incur more ambiguities.

new user-defined version is created. For example, by the command: ivi **-i bgc** : **bluo -t h~~.html** a unique THTML file of a user-defmed version **bgc 3 blue** of a web component **home**. **html** is created.

In the existing file system, both the user-defined version and its THTML files are put in the directory of their parent module. The narne convention of a THTML file of a user-defined version is **modul** *ename. encoded-user-def* **ined-versi on n-html.** The user-defined version information is encoded in encoded_user_defined_version. Variable *n* indicates the temporal version of the user-defined version. For example, **home**. *nincaoxica. 1. html* indicates the THTML source file for *english* version of **home**. html.

3.6 Conclusion and Further Directions of THTML System Design

One of the guiding principles behind the design of this THTML system was to leverage the existing web architecture and HTML language. The language **was** extended by adding TST's and TIS's. These two enhancements ailow web authors to specify which pages or components should be instantiated given a URL request from the reader. These times can be conveniently specified using relative tems like, today and tomorrow or in absolute tems using the Gregorian caiendar. Internally, these time elements are converted to six time units and comparisons between the time request and time intervals are performed using these units. Both the intensionai and extensional rules are used to resolve any ambiguities and the results of a match are sent to the client.

Since many browsers aiready exist that support conventional HTML, the system was designed to be independent of any HTML browsers. Al1 the functionality for determining which page should **be** instantiated in response to a request is located at the server. The server stores the **THTML** files, along with its time interval starnp, in a **tree** structure with a naming convention that encapsulates the parent-child relationship between the modules and the different temporal versions. Although it is less efficient than storing ail time stamps in a single file, it is more robust and facilitates implementation on a distributed file system without any bottlenecks.

The server is also responsible for interpreting requests given by the user and

generating an HTML page in response. Since users specify time using the Gregorian Calendar, a time irregularity-handling algorithm needed to **be** implemented to handle the

Figure 4.13 Three-tier Model of THTML System

irregularities when perforrning addition and subtraction operations. The server uses this algorithm to determine the correct local context from the inherited temporal context and a tag's context modifier. The temporal versioning algorithm uses this local time point to find the best-fit time interval. The generated page is sent in response to the client's request. An updated page **cm** also **be** sent at user-defined intervals using push technology. This technology leaves the responsibility with the server to keep the client up to date with the latest information. The three different tiers are shown in the **graph** above along with the communication among them. Additional tiers can also **be** added to increase the functionality of the system. Further enhancements to the system **can** include the following aspects:

At the front end, more expressive **TST's** can be added, including more userdefined temporal terminology.

At the backend, the file system **can** be separated from the web server. Due to the exponential growth of versioned entities stored in the file server, the THTML file system should be made distributed.

On the web server, another improvement would **be** to incorporate an additional

intensional tier that includes more complex temporal rules that can be used to satisfy additional requirements from web authors.

• The server push technology can be further improved by including a friendlier user **interface for web authors and more advanced protocols at backend to reduce server load.**

Chapter 5

THTML Implementation

5.1 Overview

In THTML al1 functionality is incorporated into the web server. The **web** client issues conventional **URL** and receives conventionai HïML file from THTML server. THTML is implemented independent of web **browsers.**

We witl introduce the implementation of the functionality covered in Chapter **4.** This includes the parsing of THTML files, the temporal versioning aigorithm, time irregularity handling and the generation of conventional HTML pages, dong with the data structures and the *push* implementation. They are discussed in the following order: the THTML request handler and response handler will **be** discussed **first.** Then the data structures passing between both functions are discussed. This is to keep user-defined and temporal versions. The implementation of *pusla* technology is then discussed. Finally, the chapter concludes with an example site written in THTML.

5.2 Design Requirements

THTML1.0 implements partially the designs presented in Chapter 4. With THTML1.0, TIS's cm **be** specified in an HTML file, which are the time the author intends to post the file at. The TST's can also **be** defined, which can **be** either intensional or extensional (including absolute time points) expressions. *Semer push* and the *customized calendars* are not implemented at the current stage.

In THTML 1.0, after the reader issues a request to the **web** server, a particular instance, narnely conventional HTML will **be** responded to the client at a given time. The processes at the server are expanded as follows:

Parsing the global **context**

First, the temporal context should be extracted from the URL. This context is used

by each component of the **THTML** page as the default global context.

Tirne Irregularity Handling

Time related operations are needed to retrieve the correct time context for a request. These operations are as follows:

First, the **GMT** format of **the** user-request is converted into a local time before the server can compute the global context.

Next, the local modifier and the global context are combined to retrieve the local context for each temporal element.

Figure 5.1 Enlarged THTML Web *Server*

Because of the irregularity of the Gregorian calendar, special *tirne irregulariry* algorithms are used to obtain the correct temporal context.
Instantiation

The process is illustrated in Figure **5.1.** When a temporal component is encountered in a **THTML** file, it is converted to an *extensional* representation.

First. the global context and its local context rnodifiers are combined according to the rules introduced in Section **4.2.3.1.1.** After the local context is obtained, if the component is an image or a SSI, the server will search for the corresponding THTML source and generate its instance. This instance will then **be** inserted into the newly created HTML page. The process of instantiating an instance of a temporal component is the same as converting a web cornponent with **TCS** into its extensional expression. Steps 3 to 6, above, repeats until al1 the components with **TCS** are converted. When instantiating a component, the process above is applied recursively.

5.3 THTML Implementation

5.3.1 Hook-up with APACHE API

As **was** introduced in Chapter 2, THTML is implemented by extending a module on the *Apache Semer APL*

The extended module structure is as follows:

```
modula thtml-module = € 
     STANDARD_MODULE_STUFF,
     th--ini t , /* initializar +/ 
     NULL, /* per-directory config creater */
     NULL./* dir config merger-default is to override*/
    thtml_config, /* server config creator */
     NULL, /* server config merger */
     thtml_commands, /* command table */
     th--huadlars, /* 181 lis+ of handlers +/ 
     thtml_xlate, /* [1] URI-to-filename translation */
```

```
mr /* [Il chack/vrlidata FPTP usor-id */ 
NULL. /* [5] check HTTP user id is valid *here**/
NULL, /* [3] check access by host address, etc.*/
NULL, /* [6] MIME type checker/setter */
N[fLLr /* 171 fixup~ */ 
mr /* pb80 [9] loggar / 
NtrL /+ phrsm [2J hardor prrsar */
```
The above structure module is the link between functions defined by third party and the APACHE **API.** It defines the 9 phases that the server **goes** through after receiving **a** URL request. In each phase, the third party can specify a handler to customize the server's behavior at that phase. NULL indicates that the third party is not interested in handling that request in that phase. Consequently, from the above **THTML** structure module, one can see that there **are** totdly two handlers specified to handle THTML requests, namely **thtml_handler** and **thtml_xlate,** at step [1] and [8] respectively.

};

Each handler is a function taking **as an** argument, **a** pointer to a structure calied **request-rec.** This structure contains al1 the information for a particular request, including the **user request time** in GMT format and whether the document sent to the client should be cached or not.

It also contains information to be sent to the next phase. The **thtml_xlate** handler adds user defined version information to **the** structure to **be** passed to the thtml handler.

In each phase, each handler, when invoked to handle a particular **request** rec, has to return an **int** to indicate what happened.

OK --- the request was handled successfully. This may or may not terminate the phase.

DECLINED --- no erroneous condition exists, but the module declines to handle the phase; the server tries to retrieve another.

An HTTP *exror* **code---** it aborts handling of the request,

Therefore, the request and response handlers in THTML 1.0 have the same data structure **request_rec** as input parameter and both handlers return a status code. The following sections give the functionality of them.

5.3.2 THTML Request Handler--thtml xlate

The request handler is the first phase of the server. The function is to handle the user's request by translating the **URL** to a filename.

First the URI is translated into a filename, the server will search if it is a THTML request by parsing out the version information in the filename.

If the file is a versioned entity, the server first checks whether it embeds userdefined context and temporal context. if yes, it will then parse the user-defined version information first and then the temporal context information. An error message wi11 **be** returned if it cannot **be** parsed.

The next step involves applying the best-fit algorithm to find the best-fit THTML file. The following steps are involved in this step.

The URI is translated into filename first, the filename is judged. If it is a THTML directory or a THTML file (implemented in $i_is_usexdir$), the temporal contexts and the user-defined version information if **any** are parsed out. The request time point is parsed out if there are no time contexts in the filename.(in $translate_userdir$) With the time point and the time contexts in the filename, the server finds the best fit file (in **get_time_vers**). The above steps are illustrated with Figure 5.2. In the **get_time_vers** function, the following processes are adopted to find the best-fit file: time_fit function is adopted first to judge if each time stamp in a THTML file matches the time point in the request. If yes, put the result pattern into a two dimensionai **array.** The al1 the time patterns in a file are checked, find_least_gap_time_pattern is adopted to determine the best-fit among all the candidates.

Dunng the above process, if the best-fit file cannot **be** found, the server will wnte an error message to the log file and return the error code **HTTP_NOT_POUND**.

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Otherwise in step 3, the temporal version will be resolved in two steps. First **the** time context is obtained from the THTML input, second we search for the best-fit temporal version.

Figure 5.2 Flow chart in the Request Handlet

If the target **THTML** source file is not found and the absolute THTML file does not exist, the error message **HTTP_NOT_FOUND** is given. The error is also written to a log file corresponding to the THTML file.

Before an **OK** is sent, the user-defined and temporal context information is put into the data member **user_request->config**. The response handler can use it as the global context for interpreting local temporal elements in the file.

5.3.3 THTML Response Handler- thtmi-handler

thtml-handler is the response handler, which handles the **MIME** type "text/html". It is the last function to be executed before a response is sent to the client.

Its functiondity is to translate the intensional elements in the best-fit file resolved in the request handler into extensional representations and send the resultant conventional HTML page to the client. It only handles versioned entities *SSI, image* and *hyperlink.* **~ECLINE** is sent if **the conf ig** field in the user request is **NULL.** This indicates **that** no version information **was** given.

First the server opens the best-fit source file that **was** found by the request handler

and begins to parse it. It is stored in the data structure **requst**_rec->filename. When it meets a temporal or user-defined element in the file, it is parsed into an extensional expression. In the first step, if the local context is not defined, then the global context will be used instead. It is stored in the **request** rec data structure. It performs the same procedure recursively on any included files. **First,** it checks if it **is** an **MTML** file. If yes, it will parse the version information and extract the relevant parts of the file. The function **get_time_vers** will be adopted again to retrieve the best-fit file based on a given temporal context. The above steps are illustrated **by** the following figure:

```
Whilm ( ts=I-sub-translate) 
 C 
     case ZSmCT: do-select; 
    casa ICOLLZCT: do-col 1 ect; 
     default;
1
```
in the **i-euh-tanslate** function, if **a** image, hyperlink and SSI are met, their temporal contexts will **be** parsed out, and then the best-fit will **be** found. If other tags are encountered, they are interpreted the sarne **as** the manner in conventional HTML.

```
While(next_tag) {
 if (interesting-tag) 
C 
  Fix- tag, 
  Print-tag 
1 
else if (cron tag 
       Ignore the content in it and continue parsing 
else if (include_tag)
       Do-incl ude 
else
```

```
if ( ! skep-content) 
      Print_tag
```
 \mathbf{F}

The interesting tags are the three tags that may include the temporal context switches. In the **f ix-tag** function, if a **VMOD** or **TMOD** attributes are met, their values are extracted and merged with the global intensional and temporal contexts. The resulting context is encoded into the element name. For a **SSI** and **image** tag, the best-fit file will be retrieved and included.

5.3.1. An Example

When readers go to http://valdes.uvic.ca:8080/-lliu/homepg/home.html, the following steps are performed in the **thtml-xlate** phase. The directory /-lliü/ .www/hornepg/home **.html** is checked to detennine if the directory allows global read/execute access. Then it checks if it is a directory holding multiple versions. By checking if it has a suffix . **hfml** and it is a directory, then goes to next step, check if it contains any version information. Because **home. html** carries no temporal contexts. The request time is retrieved **as** global context. Files in this directory are searched for best-fit to the request time. In this example, the file aai. I html is the best-fit. The file name dong with the user-defined contexts and temporal context are stored. Then the **the landler** will be executed. First the best-fit file will be opened, whose name is obtained in **thtml_xlate,** namely **aai.html** in this case. In this file, when the single temporal hyperlink <A **HREF="home.html" TMOD="yesterday" VMOD="lang:english"> yesterday's english version is met** first, the value of the local modifier in **VMOD** is **"langreaglish"** while there is no global version. Therefore, the archor's version wilI **be** changed to **lang r french.** The archor's name wi11 be encoded as **encoded<lang:english>6encoded<- 1~ .html** When the user clicks the above link, the **local** context will **be** the **global** context of the new page. The same steps are applied to parsing the new page.

5.4 Data Structure

The only data structure passed among the handlers is the **request** rec *structure*. This structure mainly contains protocol-related information from the client, including cache control. However, there **are** two fields that are especially relevant to **THTML,** the r equest_config and the **request_time** fields. The **request_time** is used to store the client's request time in GMT format. The **request_config** field is used to hold the structure **versionp.** This structure is where the *temporal* and *user-defined* version values are stored. The **time** versions are stored in the **tval** field while the userdefined versions are stored in the **vals** field. The data structure of *versionp* is illustrated as follows:

```
typedef struct vstruct {
            in+ ndiau; 
            char tval<sup>[50]</sup>;
            char **dims;
            char **vals:
```
1 *veraioxap;

The **dim** and **vals** variables are used to store the dimension and its comesponding version information. **ndim** indicates the number of dimensions, and **tval** is used to store the encoded time information. For exarnple, if the following user-defined version: **bgc** : **bl** *uetlang: engl* **ish+img: big+platform: unix** and temporal context:

dim	val
bgc	blue
lang	english
img	big
platform	\overline{unix}

Table 5.1 Storage of User-defined Version

1998 Dec **I2/11:11:13** are encountered, **they** will **be** stored in *versionp* While *tval*

stores

Y1998M12D12H11I11S13 (The format and meaning will be discussed shortly). **ndim** is 4 indicating that there are four user-defined dimensions.

This structure hoids the global tirne and user-defined context, which will **be** used in the response handler to interpret the local intensionai **and** temporal elernents.

5.5 Retrieving the best-fit THTML file

S.S. 1 Retrieving Candidates

Function **get_time_vers** is used to find the best-fit interval candidate. The process of **finding** candidates is the process of finding if the time point in a **query** lies in the time interval starnps defined in the THTML **files.**

In order to compare them, THTML puts the parsed time pattern found between the **<cron>** and **</cron>** tags into the following structure.

```
typedœf struct th{
```

```
char<sup>*</sup> sec;
                char* zain; 
                char* hour ; 
                char* dom; 
                char* month;
                char* yorr; 
                char* dw; 
                int flags
>tiauS;
```
The fields hold respectively: second, minute, hour, day of month, month, year and day of **week.** The last **flag** indicates the time pattern, which can **be** one of the constants **BEFORE, AFTER, DURING, RECURRING.** THTML encodes the two end points of a time unit into one field and **decodes** it later when cornparhg it with a time point. Since

only 2 digits are needed to represent each time unit, except for the year, then we can encode the upper endpoint for each unit as the 3rd and 4th and the lower endpoint in the 1st and 2nd position. The year is encoded in a sirnilar manner except using 4 digits each for the upper **and** lower endpoint. When comparing times, we always start with the most significant unit, so we compare the year first, followed by *month*, date, hour, minute and second. Unnecessary cornparisons are avoided by discarding elements as soon **as** one time unit does not match.

There are several different situations that need to be handled by the server. These situations will be described below.

• BEFORE/AFTER pattern comparison:

For the **BEFORE** pattern, when we compare each time unit, **the** function **get_time_vers** returns a 1 if the time request is less than the time stamp, $a - 1$ if it is greater than the time starnp or continue with the next unit if they are equal. If the **2** time points are exactly equal, we also return **an** error, namely -1. **A** similar situation oçcurs for **AFTER** patterns.

If a **DffRING** pattem is encountered, the corresponding unit will **be** decoded first.

When the first data type is encountered in one time stamp that cannot hold the data in timestamp exclusively, the matching will stop and the function **get_time_vers** returns -1 to the calling function indicating failure; otherwise, if there is **any** time unit that contains the corresponding data unit in time point exclusively, the function retums immediately with the value **1.** Because no matter what value is in the remaining smdler data units, it is always fits in **the** time pattern; **Lf** either of them overlaps with it, the process of comparison will continue recursively with the remaining time units until it is clear whether the time point fits in the tirne pattem or not. At the end, it tums out that the tying can not be solved between the two time points, *i.e.* the time point is overlapped with one of the end point of the timestamp, 1 is returned.

If a recurring pattem is encountered, the server compares each individual time unit. If a time unit containing intervais is encountered, it will **be** treated in the same manner as a 'during' pattern. The one thing that distinguishes **this** pattern from al1 others

is that if a match is found in a time unit, only the matched value is stored and the other

BEFORE:

Figure 5.3 Algorithm of Refinement Relationship Resolution between "before" Patterns

recurring values are discarded.

The situation is simplest in a comparison between two time points. In this case, we just need to match al1 time units.

If a match is found, the candidate is stored in an **array** so that it **can be** compared later with other candidates to determine the best-fit.

5.5.2 Resolving Ambiguity

In the following section, only ambiguities among temporal versions are discussed. The ambiguity of user-defined versions is detailed in the corresponding literature [10].

The ambiguity resolution is handled in the function find_least_gap_time_pattern, which get_time_vers uses to find the best-fit THTML file or the best-fit time interval in a single file. In this function, if more than one pattern matches the time request, the rules introduced in Chapter 3 are applied to solve the ambiguity. The matched patterns are stored in a two-dimensional **may.** The

Table 5.2 Temporal Pattern Candidates

structure is illustrated in Table 5.2.

The comparison returns the best-fit pattern for further comparison among the candidates, or -1 to the calling function **get_time_vers** indicating there is ambiguity, in which case the server finds the absolute **THTML** file instead as the best fit. If it does not exist, a log is written in THTML log file indicating the vanilla version does not exist.

The server **goes** through al1 the candidates, to find the pattem with the highest refinement level. (The order of refinement, from highest to Iowest, is *point, during, after/before, and recurring.)*

If only one pattern is found in the server, it will return that pattern.

Note that if one of the matches is a *point* pattern, it **is** the only one (since time points are unique). So it becomes our best-fit.

• If there is more than one *during* pattern, the difftime function is called to calculate the time intervaI. The smallest interval is the best fit.

If both a *before* and an *a/er* pattern exist the most refined pattern, the error code indicating ambiguity is returned in this situation because there is no way to compare the time gap of these two patterns.

If the most refinement patterns are either of *afler* or **before, the** value of *dflime* between the end point and the target end point is the criterion to choose the best-fit candidate. The one having smaller value is the best-fit indicating that it encapsulates more closely the time point requested.

If there is more than one *recwring* pattem candidate, the algorithm to find the best-fit is similar to the interval pattern, the time units are compared from the highest to the lowest.

The ambiguity algorithm among files or in one file is the same, except that in the ambiguity handling among files, a file index is kept dong **with** its best-fit pattern in the data structure. When the ambiguity is solved, the corresponding file index is returned so that in the response handler, the server can open the best-fit file.

5.6 Time irregularity Handling

As mentioned in Chapter 4, special functions (drivers) are provided to handle time irregularity.

The functions provided are based on the Gregorian Calendar:

1. getDMonth, which returns the number of days in a given month and year.

2. getDow, which returns the day of week for the given date.

When **THTML** merges local and global temporal contexts, THTML merges the **two** time patterns starting from the smallest time unit to the Iargest, i.e. **from** the seconds to the years.

In **THTML,** the day of the **week** information is calculated when necessary and so it stores the date using only the 6 time units mentioned above. The reason for this is as follows:

1. to simplify the algorithm

2. **to** avoid mismatches in the date and the day of week given. For example, if the user gives the time pattern:

f + **5 4 1996 1**

and April 5 happens to not be a Monday. THTML assumes that the date is correct in this situation rather than the day **of** week. There are however situations when we need to **convert** from the date to a day **of** the week and vice-versa.

When combining local and global contexts, THTML has to convert the week information given by a user into a date. For example, when a user defines **TMOD="1d+2w"** in his query, it adds $2*14+1$ days to the global context.

Another situation is when it meets a timestamp indicating the **dorv.** THTML has to calculate the *dow* of the time point in clients' request for comparison.

For example, the timestamp

1 2 3 5 1998 1,2,3

is compared with the time point **May IZ,12:12:I2,1998.** The **THTML** server will calculate the day of week of the time point to check if it matches the **dow** given by the timestamp.

Figure 5.4 "Overflow" Among Time Units

5.7 Other Aspects in Implementation

Version Encoding and Version Decoding

As seen in the design section, THTML encodes the version information in the file name and decodes it when the server retrieves the best-fit version. There are a couple of reasons why it is encoded. The main reason is to prevent the laymen from hacking the system **based** on the version information. The second reason is to standardize the version information (especially user-defined versions) with standard printable ASCII characters.

In THTML, the *absolute* and relative time contexts are encoded using different patterns so that the server can easily differentiate between the two and process them accordingly.

For example, if a link

is encountered, the time context is written as **SOMOIOD-13MOYO** and encoded

foIIowing the filename of the link with a separating symbol **Q** in between.

When the user clicks on it, the server will decode the time information. Since it is relative tirne, a global context is needed for it to merge with. Otherwise the globai time context is ignored and the decoded absolute time is used instead as the time for the query.

Implementation of *Server* **Push**

Server push was not implemented in the current THTML software for the following reason:

In order to implement server push technology, we need the ability to send multiple parts of a single document using the **HTïP/1.0** protocol. However **APACHE-1.2.5** has a bug where the server cannot **be** modified to keep multiple connections open and efficiently close them at the user's request. The details of the bug can be found in *APACHE'S* bug reports.

5.8 An Example Site

In order to clarify the design in THTML server. a simple example site is designed and implemented. The experiment site is about the horoscopes. The time of clients' request determines the corresponding page to **be shown.** Each of them also have versions having the picture of the sign, or not, the user **can** also change the background of the page according to their preference.

Each **page** is linked with the previous sign page and the next sign page, so that the client can iteraie al1 of the sign pages from either of them. The page combines temporal and user-defined HTML together. One can see **how** THTML creates a *dynamic* sire. With conventional **HTML,** if each page has three versions of pictures and three versions of background, the total version will be $3*3$ pages, each has the storage equal to the combined ones of the background and picture version.

While in THTML, $3+3$ files are needed for the same amount of versions of the same site. The storage of the file is the same as each of them, as seen above, THTML Save

huge storage in versioned site compared with conventionai HTML.

In the above THTML, the client wilI see different pages of horoscope depending on the time of retneving the page. The time accords to the time of horoscope at that time. In it, the request of previous and next one are also implemented with requests for **particular** temporal instances. The following two screenshots are a **THTML** source file and the corresponding conventional HTML page in response at a particular time. **Fiam** this **example,** one can see how to create a temporal site in **THTML.**

```
<HTML>
<HEAD>
          <TITLE>Capricom information - All about Capricom!</TITLE>
<!-- Changed by:, 13-Jan-2000 -->
<fHEAD>
<BODY background="Capricorn information - All about Capricorn! files/whtmarb.gif">
<CENTER><IMG alt="capricorn" border="0" height="121"
src="Capricom information - All about Capricom! files/capricom.gif" width="161">
<CENTER>
H1<CENTER><BR>The Goat<BR>December 22 to January 20</CENTER></H1>
<FONT color="#008000" size="+2"><B>
<HR><TABLE border="1" cellPadding="3" cellSpacing="3">
 <TBODY>
 <TR>
<TH align="middle" vAlign="top">
<A href="home.M1DnCgOxjcq@JaaervXx9eh-ebwm3F9VSA_6FXEwgae.html">
<IMG border="0" src="Capricorn information - All about Capricorn!_files/button2.gif">
<BR> Last Sign
          </A> </TH>
  <TH align="middle" vAlign="top">
<A href="home M1iDnCgOxicq@JaservXzJOXDQDaF7VLs4jFRnIlai.html">
<IMG border="0" src="Capricom information - All about Capricom! files/button2.gif">
<BR> Next Sign
  </A></TH>
</U></TR>
</TABLE></FONT>
```
<BODY><HTML>

Figure 5.5 An Extension Sent to a Client at the Request Tirne 10:10:58 am, Jan 14,OO

E home MillinLqUxicqi | Notepad

تتر د

st File Edit Search Help <cmp> *** 22-20 12-1 * <kmp> <HTML> <HEAD> <TITLE>Capricorn information - All about Capricorn!</TITLE> <!-- Changed by: , 13-Jan-2000 --> </HEAD> <BODY background="Capricorn information - All about Capricorn! files/whitmarb.gil"> <CENTER> $<$ CENTER> $<$ H₁> <CENTER>
The Goat
December 22 to January 20<CENTER></H1> H <TABLE border="1" cellPadding="3" cellSpacing="3"> <TBODY> $<$ TR> <TH align="middle" vAhgn="top">
 Last Sign </TH>

```
<TH align="middle" vAlign="top">
```


 Next Sign </TH>

```
</U></TR>
</TABLE></FONT>
```
</BODY></HTML>

Figure 5.6 The THTML Source for the Generated HTML Page Shown in Figure 5.5

Chapter 6

Conclusion and Further work

6.1 THTiML----Dynamic HTML

There are two ways in which HTML pages can be dynamic. One method is to have the client's browser dynamically interpret an html file. **DHTML** belongs to this category. In DHTML, dynamic features can be implemented with functions written in Javascript that can **be** interpreted and run by the clients' browsers. The other method is to have the server continually generate **HTML** files on the fly and sending them to the client. Examples of this method include Microsoft's ASP and **XML.** THTML introduced in this thesis also belongs to this category.

A THTML server behaves differently **than** an HTTML server.

First, it provides temporal instances of a URL at a particular time context. It is able to do this by first, providing a system to store and retrieve the most relevant timestamped HTML sources. Second, it also allows the client to ernbed time sensitive tags into conventional HTML components. Third, with push technology incorporated, THTML sites are updated using user-defined frequencies and preferences.

The above features are based on two major aspects in design. One is the TIS design in a temporal HTML file; the other is the refinement ordering of the time intervals.

In the former, the TIS is the time that a web document is expected to post. THTML combines the two data types in its representation: *date* and *time* to cover the required ganularity. The format is as follows:

second minute hour **day** *month* **year [dow]**

With the above data type, basic time patterns in **web** authors' demands can **be** represented, which includes *before, afier, during, time point* and *recurring* pattern.

In the latter, the extensional and intensional rules **are** specified in order to minimize

the ambiguity that may occur if multiple time intervals satisfy **the** tirne point in a client's request. The extensional mle specifies that the smaller **tirne** gap of the time interval stamp to a particular time point, the better-fit of the THTML source file it associates with. The intensionai rule further refines the extensional rule to help resolve any ambiguities. The refinement order of the five time patterns that THTML supports is (from the lowest to the highest) **recirrring, before/after,** *during. tirne* **point.** Among them, *before* and *after* have the same priority. Consequently, refinement ordering is not only related to the temporal gap but also the type of intervals. With the above algorithm, the ambiguities involved in finding the best-fit interval can **be** minimized. Therefore, the instance **as** server's response is generated using the THTML file whose time interval stamp fits the time point in request best.

6.2 Further Implemen tations

THTML was designed to fulfill more complex temporal requirements at the frontend. However some of these features still need to **be** implemented. They are:

Server push technology needs to **be** integrated into the system. The server push interface needs to **be** enhanced with more options like alIowing an aiarm and pop up window when the site is updated. It should also allow the user to define the frequency of the updates.

The file system can be isolated from the file server **and** developed into a large file pool. If temporal files are perceived as temporal data, ail transactions can be perceived as storing data in one big repository and allowing the queries to retrieve the needed data from this repository.

• Natural language can be used in place of the time sensitive tags to make it more user friendly. In order to support the new features at the front end, another intensional tier can be incorporated at the back-end. This may include compiex reductive niles and mathematical expressions to support more complex temporal requirements at the frontend. For example, if the author intends to post a page with a complex time pattern according to the Fibonacci function based on the two initial time points. With the Fibonacci rules incorporated at the backend, THTML is expected to fulfill the

requirement.

6.3 Possibility of Combining THTML, DHTML and XML

Although the three enhancements of HIML provide *dynamic* solutions from different perspectives, due to the fact that they are **al1** based on hierarchical web document philosophy, there is possibility of combining them in use.

With Java **and** JavaScript

lava is an object-oriented programming language that can be run on Web. It is embedded in HTML file, the browser finds the Java class file in the server and runs the program with Java Virtual machine instailed.

Javascript is an object-oriented scripting language embedded in HTML file. By defining

RVNAT tag, the browser **can** run the script at client side or server side. One enhancement is making it more interactive with callback functions.

DHTML perceives **any** part of the content as an object, and allows interaction and behavior of each with JavaScript and **CCS.**

THTML is the versioning soiution of a web site with time With DHTML combined, any part of the web content cm **be** distinct as a temporal versioned object (web elements). It will make the definition of web document beyond the limitation of time and content, which will give the web author more flexibility in designing web documents.

A Java and the object including server side Javascript (Serverlet) can also **be** made temporal versioned entities. Based on different time, different Java program will **be** triggered at server side. (The author has no clue of how to combine the client-site JavaScript with THTML)

With *XML*

XML allows user-defined **tags** in HTML file. There is a **XML** parser to parse it. Based on the particular data in particular user-defined tag, particular business logic is adopted at web server. From this aspect, it is similar to THTML. If THTML is cornbined, XML elements **can** have multiple versions in time, different versions can be stored at server side. By combining the two, the user can write a document with user**defined** styles as well as user-defined frequencies.

6.4 Prospect

One can see that as tirne-sensitive information on the **web** becornes more prevalent, the greater the need for adoption of THTML for such applications as newspaper and stock market sites. **It** ailows the web readers to retrieve a version with a **particular** time. **It** also facilitates the planning of a site that changes very frequently without the user having to update it with the same frequency. With further implementation and the combination with the new **web** technologies, **THTML** will become a powerful authoring tool for *dynamic* web sites in advanced **web** applications.

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